

FILL THE BLANK, A TREE OF N VERTICES HAS _____ EDGES. GROUP OF ANSWER CHOICES $N-1$ N Nn $2N$

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UNDERSTANDING THE BASICS: WHAT IS A TREE IN GRAPH THEORY?

BEFORE DIVING INTO THE SPECIFICS OF THE NUMBER OF EDGES IN A TREE, IT'S ESSENTIAL TO UNDERSTAND WHAT A TREE IS WITHIN THE CONTEXT OF GRAPH THEORY. GRAPH THEORY IS A BRANCH OF MATHEMATICS AND COMPUTER SCIENCE THAT STUDIES GRAPHS, WHICH ARE STRUCTURES MADE UP OF VERTICES (ALSO CALLED NODES) CONNECTED BY EDGES (ALSO CALLED LINKS).

A TREE IS A SPECIAL TYPE OF GRAPH CHARACTERIZED BY THE FOLLOWING PROPERTIES:

- IT IS CONNECTED: THERE IS A PATH BETWEEN ANY TWO VERTICES.
- IT IS ACYCLIC: IT CONTAINS NO CYCLES OR LOOPS.
- IT HAS N VERTICES (OR NODES): THE NUMBER OF POINTS IN THE GRAPH.

KEY DEFINITIONS

VERTICES (NODES): THE FUNDAMENTAL UNITS OR POINTS IN A GRAPH.

EDGES (LINKS): THE CONNECTIONS BETWEEN VERTICES.

VISUAL REPRESENTATION OF A TREE

IMAGINE A FAMILY TREE OR A COMPANY'S ORGANIZATIONAL CHART. THESE ARE REAL-WORLD EXAMPLES THAT RESEMBLE THE STRUCTURE OF A TREE IN GRAPH THEORY—CONNECTED AND WITHOUT CYCLES.

THE CORE QUESTION: HOW MANY EDGES DOES A TREE HAVE?

THE MAIN QUESTION ADDRESSED IN THIS ARTICLE IS: "A TREE OF N VERTICES HAS HOW MANY EDGES?"

THIS QUESTION IS FUNDAMENTAL IN UNDERSTANDING THE STRUCTURE OF TREES, WHICH ARE PREVALENT IN VARIOUS APPLICATIONS SUCH AS NETWORK DESIGN, DATA STRUCTURES, AND ALGORITHMS.

THE FILL-IN-THE-BLANK STATEMENT

THE STATEMENT TO ANALYZE IS:

"FILL THE BLANK, A TREE OF N VERTICES HAS _____ EDGES. GROUP OF ANSWER CHOICES $N-1$ N Nn $2N$ "

THE OPTIONS ARE:

- $N-1$
- N
- Nn
- $2N$

OUR GOAL IS TO DETERMINE THE CORRECT CHOICE AND UNDERSTAND WHY IT IS SO.

FORMAL EXPLANATION: THE MATHEMATICAL RELATIONSHIP IN TREES

THEOREM: NUMBER OF EDGES IN A TREE

IN GRAPH THEORY, THERE IS A WELL-KNOWN THEOREM:

> "A TREE WITH N VERTICES ALWAYS HAS EXACTLY $N - 1$ EDGES."

THIS STATEMENT IS FUNDAMENTAL AND HOLDS TRUE FOR ALL TREES, REGARDLESS OF THEIR SIZE OR COMPLEXITY.

WHY IS THIS TRUE?

THE INTUITIVE EXPLANATION HINGES ON THE PROPERTIES OF A CONNECTED, ACYCLIC GRAPH:

- TO CONNECT N VERTICES IN A TREE WITHOUT FORMING CYCLES, YOU NEED TO CONNECT EACH NEW VERTEX TO THE EXISTING STRUCTURE WITH EXACTLY ONE EDGE.
- STARTING FROM A SINGLE VERTEX, EVERY NEW VERTEX ADDED REQUIRES A SINGLE EDGE TO CONNECT IT TO THE EXISTING STRUCTURE.
- THEREFORE, FOR N VERTICES, THE TOTAL NUMBER OF EDGES IS $N - 1$.

FORMAL PROOF SKETCH

BASE CASE: FOR $N = 1$ (A SINGLE VERTEX), THE NUMBER OF EDGES IS 0 , WHICH ALIGNS WITH $N - 1$.

INDUCTIVE STEP:

- ASSUME A TREE WITH N VERTICES HAS $N - 1$ EDGES.
- WHEN ADDING A NEW VERTEX, TO KEEP THE GRAPH CONNECTED AND ACYCLIC, EXACTLY ONE NEW EDGE IS NEEDED.
- THUS, FOR $N + 1$ VERTICES, THE TOTAL EDGES ARE $(N - 1) + 1 = N$, WHICH IS $(N + 1) - 1$.

THIS INDUCTION CONFIRMS THAT THE NUMBER OF EDGES IN A TREE WITH N VERTICES IS ALWAYS $N - 1$.

EXPLORING THE ANSWER CHOICES

LET'S ANALYZE EACH OF THE OPTIONS PROVIDED:

1. $N - 1$

- CORRECT ANSWER: YES.
- EXPLANATION: AS ESTABLISHED, A TREE WITH N VERTICES HAS EXACTLY $N - 1$ EDGES.

2. N

- INCORRECT: A TREE WITH N VERTICES DOES NOT HAVE N EDGES; THAT WOULD IMPLY AN EXTRA CONNECTION, POTENTIALLY CREATING CYCLES.

3. N^N

- INCORRECT: THIS SUGGESTS THE NUMBER OF EDGES DEPENDS ON THE PRODUCT OF N AND SOME OTHER VARIABLE n , WHICH ISN'T RELEVANT HERE UNLESS n IS SPECIFIED. THE STANDARD FORMULA DOESN'T INVOLVE SUCH A MULTIPLICATION.

4. 2^N

- INCORRECT: SIMILAR TO THE PREVIOUS, THIS DOES NOT ALIGN WITH THE FUNDAMENTAL PROPERTY OF TREES.

APPLICATIONS OF THE $N - 1$ EDGES PROPERTY IN REAL-WORLD SCENARIOS

UNDERSTANDING THAT A TREE WITH N VERTICES HAS $N - 1$ EDGES ISN'T JUST A THEORETICAL CONSTRUCT; IT HAS PRACTICAL IMPLICATIONS.

NETWORK DESIGN

- WHEN DESIGNING A MINIMAL NETWORK THAT CONNECTS N COMPUTERS OR DEVICES WITHOUT REDUNDANCY, YOU NEED AT LEAST $N - 1$ CONNECTIONS.
- ADDING MORE CONNECTIONS INTRODUCES CYCLES, WHICH MAY BE UNNECESSARY AND COSTLY.

DATA STRUCTURES

- BINARY TREES: A COMMON DATA STRUCTURE IN COMPUTER SCIENCE, WHERE THE NUMBER OF EDGES (OR LINKS) ALSO FOLLOWS THE $N - 1$ RULE IN THE CASE OF SIMPLE TREES.
- MINIMUM SPANNING TREES: ALGORITHMS LIKE KRUSKAL'S OR PRIM'S FIND THE SUBSET OF EDGES CONNECTING ALL VERTICES WITH THE MINIMAL TOTAL WEIGHT, ALWAYS RESULTING IN $N - 1$ EDGES.

ORGANIZATIONAL HIERARCHIES

- IN AN ORGANIZATIONAL CHART MODELED AS A TREE, THE NUMBER OF REPORTING LINES (EDGES) CONNECTING N EMPLOYEES (VERTICES) IS $N - 1$.

VISUALIZING THE RELATIONSHIP: EXAMPLES

LET'S CONSIDER SMALL VALUES OF N AND VISUALIZE THE NUMBER OF EDGES:

NUMBER OF VERTICES (N)	NUMBER OF EDGES ($N - 1$)	EXAMPLE STRUCTURE
1	0	SINGLE NODE WITH NO CONNECTIONS
2	1	TWO NODES CONNECTED BY A SINGLE EDGE
3	2	A CHAIN OR A STAR WITH 3 NODES
4	3	A TREE WITH 4 NODES, LIKE A LINE OR A STAR

THESE EXAMPLES REINFORCE THE $N - 1$ RULE.

COMMON MISCONCEPTIONS AND CLARIFICATIONS

IS $N - 1$ ALWAYS TRUE?

- YES, FOR TREES SPECIFICALLY. OTHER GRAPHS MAY HAVE DIFFERENT RELATIONSHIPS:
- A COMPLETE GRAPH WITH N VERTICES HAS $N(N - 1)/2$ EDGES.
- A CYCLE GRAPH WITH N VERTICES HAS N EDGES.

DOES ADDING EDGES CREATE CYCLES?

- YES, ADDING ANY EXTRA EDGE BEYOND $N - 1$ IN A TREE CREATES AT LEAST ONE CYCLE, VIOLATING THE ACYCLIC PROPERTY OF A TREE.

SUMMARY AND FINAL THOUGHTS

- THE FUNDAMENTAL PROPERTY OF TREES IN GRAPH THEORY STATES THAT A TREE WITH N VERTICES ALWAYS HAS $N - 1$ EDGES.
- THIS PROPERTY IS CRITICAL IN DESIGNING EFFICIENT NETWORKS, UNDERSTANDING HIERARCHICAL STRUCTURES, AND DEVELOPING

ALGORITHMS.

- RECOGNIZING THIS RELATIONSHIP HELPS IN SOLVING MANY PROBLEMS RELATED TO CONNECTIVITY, OPTIMIZATION, AND GRAPH TRAVERSAL.

CONCLUSION: THE CORRECT CHOICE

BASED ON THE DETAILED ANALYSIS, THE CORRECT ANSWER TO THE FILL-IN-THE-BLANK QUESTION IS:

$N - 1$

A TREE OF N VERTICES HAS $N - 1$ EDGES.

ADDITIONAL RESOURCES FOR FURTHER LEARNING

- GRAPH THEORY TEXTBOOKS: EXPLORE MORE ABOUT TREES, CYCLES, AND CONNECTIVITY.
- ALGORITHMS: STUDY ALGORITHMS LIKE KRUSKAL'S AND PRIM'S FOR MINIMUM SPANNING TREES.
- ONLINE COURSES: PLATFORMS LIKE COURSERA AND EDX OFFER COURSES ON GRAPH THEORY AND NETWORK DESIGN.

FAQs

Q1: CAN A TREE HAVE MORE THAN $N - 1$ EDGES?

A1: NO. ADDING MORE EDGES RESULTS IN CYCLES, MAKING IT NO LONGER A TREE.

Q2: WHAT HAPPENS IF A TREE HAS FEWER THAN $N - 1$ EDGES?

A2: THE GRAPH WOULD BE DISCONNECTED, MEANING IT ISN'T A TREE.

Q3: ARE THERE SPECIAL TYPES OF TREES WITH DIFFERENT PROPERTIES?

A3: YES. FOR EXAMPLE, A FULL BINARY TREE, AVL TREE, OR RED-BLACK TREE EACH HAS UNIQUE PROPERTIES, BUT THE $N - 1$ EDGES RULE STILL APPLIES FOR BASIC TREES.

BY UNDERSTANDING THE RELATIONSHIP BETWEEN VERTICES AND EDGES IN TREES, YOU CAN BETTER GRASP GRAPH STRUCTURES AND THEIR APPLICATIONS ACROSS VARIOUS FIELDS.

FREQUENTLY ASKED QUESTIONS

IN A TREE WITH N VERTICES, HOW MANY EDGES DOES IT HAVE?

$N - 1$

WHAT IS THE NUMBER OF EDGES IN A TREE CONSISTING OF N VERTICES?

$N - 1$

FOR A TREE WITH N VERTICES, WHICH OF THE FOLLOWING CORRECTLY REPRESENTS THE NUMBER OF EDGES?

$N - 1$

IN GRAPH THEORY, A TREE WITH N VERTICES ALWAYS HAS HOW MANY EDGES?

$N - 1$

CHOOSE THE CORRECT NUMBER OF EDGES FOR A TREE WITH N VERTICES FROM THE OPTIONS: $N-1$, N , Nn , $2n$.

$N - 1$

ADDITIONAL RESOURCES

FILL THE BLANK, A TREE OF N VERTICES HAS _____ EDGES IS A FUNDAMENTAL QUESTION IN GRAPH THEORY THAT OFTEN APPEARS IN MATHEMATICS, COMPUTER SCIENCE, AND COMBINATORICS. UNDERSTANDING THE RELATIONSHIP BETWEEN THE NUMBER OF VERTICES AND EDGES IN A TREE IS CRUCIAL FOR ANALYZING THE STRUCTURE AND PROPERTIES OF VARIOUS GRAPH TYPES. THIS ARTICLE PROVIDES A COMPREHENSIVE GUIDE TO THIS CONCEPT, EXPLORING THE DEFINITION OF A TREE, ITS CHARACTERISTICS, AND THE REASONING BEHIND THE SPECIFIC NUMBER OF EDGES IT CONTAINS, CULMINATING IN AN IN-DEPTH EXPLANATION OF WHY THE CORRECT ANSWER IS $N - 1$.

UNDERSTANDING THE BASICS: WHAT IS A TREE?

DEFINITION OF A TREE

A TREE IS A TYPE OF GRAPH THAT IS BOTH CONNECTED AND ACYCLIC. THIS MEANS:

- CONNECTED: THERE IS A PATH BETWEEN ANY TWO VERTICES IN THE GRAPH.
- ACYCLIC: THE GRAPH CONTAINS NO CYCLES OR CLOSED LOOPS.

KEY PROPERTIES OF A TREE

- NUMBER OF EDGES: A TREE WITH N VERTICES HAS EXACTLY $N - 1$ EDGES.
- MINIMAL CONNECTIVITY: A TREE IS THE MINIMAL NUMBER OF EDGES NEEDED TO KEEP THE GRAPH CONNECTED.
- UNIQUE PATH: BETWEEN ANY TWO VERTICES, THERE EXISTS EXACTLY ONE SIMPLE PATH.

THESE PROPERTIES MAKE TREES A FUNDAMENTAL STRUCTURE IN VARIOUS APPLICATIONS, SUCH AS DATA STRUCTURES (BINARY TREES, SPANNING TREES), NETWORK DESIGN, AND HIERARCHICAL MODELING.

THE RELATIONSHIP BETWEEN VERTICES AND EDGES IN A TREE

WHY DOES A TREE HAVE $N - 1$ EDGES?

THE CORE REASON IS ROOTED IN THE MINIMAL CONNECTIVITY REQUIREMENT. TO UNDERSTAND THIS, CONSIDER THE FOLLOWING:

- FOR A GRAPH WITH N VERTICES TO BE CONNECTED, IT MUST AT LEAST HAVE ENOUGH EDGES TO CONNECT ALL VERTICES WITHOUT FORMING CYCLES.
- REMOVING ANY EDGE FROM A TREE WOULD DISCONNECT THE GRAPH, INDICATING THAT ALL EDGES ARE CRITICAL FOR MAINTAINING CONNECTIVITY.
- CONVERSELY, ADDING ANY ADDITIONAL EDGE WOULD CREATE A CYCLE, VIOLATING THE ACYCLIC PROPERTY OF A TREE.

FORMAL PROOF SKETCH

TO SOLIDIFY THIS CONCEPT, HERE IS A SIMPLIFIED PROOF:

1. BASE CASE: A TREE WITH ONE VERTEX ($N=1$) HAS ZERO EDGES. SO, WHEN $N=1$, $\text{EDGES} = 0 = 1 - 1$.

2. INDUCTIVE STEP: ASSUME THAT ANY TREE WITH N VERTICES HAS $N - 1$ EDGES.

3. ADDING A VERTEX: TO ADD A NEW VERTEX TO A TREE WITH N VERTICES, CONNECT IT WITH EXACTLY ONE NEW EDGE TO AN EXISTING VERTEX (TO KEEP THE GRAPH CONNECTED AND ACYCLIC).

4. RESULT: THE NEW TREE HAS $N + 1$ VERTICES AND $(N - 1) + 1 = N$ EDGES, MAINTAINING THE PROPERTY.

THIS INDUCTIVE REASONING CONFIRMS THAT A TREE WITH N VERTICES HAS $N - 1$ EDGES.

EXPLORING THE MULTIPLE CHOICE OPTIONS

GIVEN THE QUESTION: FILL THE BLANK, A TREE OF N VERTICES HAS _____ EDGES, THE CHOICES ARE:

- $N - 1$
- N
- N^N
- 2^N

LET'S ANALYZE EACH OPTION IN DETAIL.

1. $N - 1$

- CORRECT CHOICE. AS ESTABLISHED, A TREE WITH N VERTICES ALWAYS HAS $N - 1$ EDGES.
- WHY IT FITS: THIS IS A FUNDAMENTAL THEOREM IN GRAPH THEORY, OFTEN PROVEN VIA INDUCTION OR COMBINATORIAL ARGUMENTS.

2. N

- INCORRECT. A TREE WITH N VERTICES DOES NOT HAVE N EDGES; IT IS ONE LESS.
- COMMON MISCONCEPTION: SOMETIMES PEOPLE CONFUSE THE TOTAL NUMBER OF VERTICES WITH THE NUMBER OF EDGES, BUT IN TREES, THE EDGES ARE ALWAYS ONE FEWER THAN VERTICES.

3. N^N

- INCORRECT. THIS SUGGESTS A MULTIPLICATIVE RELATIONSHIP THAT DOES NOT APPLY TO TREES.
- CONTEXT: THIS MIGHT BE RELEVANT IF CONSIDERING COMPLETE BIPARTITE GRAPHS OR OTHER STRUCTURES, BUT NOT TREES.

4. 2^N

- INCORRECT. SIMILAR TO THE PREVIOUS POINT, THIS IS NOT A CHARACTERISTIC OF TREES.
- NOTE: FOR CERTAIN GRAPHS LIKE BIPARTITE GRAPHS, THIS COULD BE RELEVANT, BUT NOT FOR TREES.

PRACTICAL IMPLICATIONS AND APPLICATIONS

UNDERSTANDING THE RELATIONSHIP BETWEEN VERTICES AND EDGES IN TREES HAS PRACTICAL SIGNIFICANCE:

- NETWORK DESIGN: ENSURING MINIMAL CONNECTIONS FOR MAXIMUM REACHABILITY.
- DATA STRUCTURES: BINARY TREES AND SPANNING TREES RELY ON THIS PROPERTY.
- ALGORITHM OPTIMIZATION: EFFICIENT ALGORITHMS FOR TRAVERSAL, SEARCH, AND PATHFINDING OFTEN LEVERAGE TREE PROPERTIES.

EXAMPLES IN REAL LIFE

- COMPUTER NETWORKS: DESIGNING A NETWORK WITH N DEVICES CONNECTED WITH $N - 1$ LINKS MINIMIZES COSTS WHILE MAINTAINING CONNECTIVITY.

- HIERARCHICAL STRUCTURES: ORGANIZATIONAL CHARTS OR FILE DIRECTORY TREES ARE MODELED AS TREES, WHERE EACH NODE (DEPARTMENT OR FOLDER) IS CONNECTED VIA $N - 1$ LINKS.

SUMMARY: THE CORRECT ANSWER

IN CONCLUSION, THE NUMBER OF EDGES IN A TREE OF N VERTICES IS $N - 1$. THIS FUNDAMENTAL FACT UNDERPINS MANY CONCEPTS IN GRAPH THEORY AND COMPUTER SCIENCE.

FINAL ANSWER: $N - 1$

FILL THE BLANK, A TREE OF N VERTICES HAS $N - 1$ EDGES.

ADDITIONAL NOTES

- SPECIAL CASES: FOR $N = 1$, THE TREE HAS 0 EDGES.
- CYCLE DETECTION: IF A GRAPH WITH N VERTICES HAS MORE THAN $N - 1$ EDGES, IT MUST CONTAIN AT LEAST ONE CYCLE.
- SPANNING TREES: IN WEIGHTED GRAPHS, A SPANNING TREE IS A SUBSET OF EDGES FORMING A TREE THAT INCLUDES ALL VERTICES WITH $N - 1$ EDGES, OFTEN USED IN ALGORITHMS LIKE KRUSKAL'S AND PRIM'S.

BY GRASPING THIS CORE CONCEPT, YOU'LL HAVE A SOLID FOUNDATION FOR UNDERSTANDING MORE COMPLEX TOPICS IN GRAPH THEORY, NETWORK DESIGN, AND ALGORITHM DEVELOPMENT.

Fill The Blank A Tree Of N Vertices Has Edges Group Of Answer Choices N 1 N N N 2 N

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