

YOU ARE GIVEN A STRING S CONSISTING OF N LOWERCASE LETTERS OF THE ENGLISH ALPHABET. FIND THE LENGTH OF

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INTRODUCTION TO STRING LENGTH CALCULATION

IN THE REALM OF COMPUTER SCIENCE AND PROGRAMMING, STRINGS ARE FUNDAMENTAL DATA TYPES USED TO REPRESENT TEXTUAL DATA. A STRING IS A SEQUENCE OF CHARACTERS, AND UNDERSTANDING HOW TO DETERMINE ITS LENGTH—THAT IS, THE NUMBER OF CHARACTERS IT CONTAINS—IS A BASIC YET CRUCIAL OPERATION. WHETHER YOU'RE DEVELOPING APPLICATIONS, SOLVING ALGORITHMIC PROBLEMS, OR WORKING ON DATA ANALYSIS, CALCULATING THE LENGTH OF A STRING IS OFTEN THE FIRST STEP IN STRING MANIPULATION TASKS.

IN THIS ARTICLE, WE WILL EXPLORE THE PROBLEM STATEMENT: GIVEN A STRING S CONSISTING OF N LOWERCASE ENGLISH LETTERS, HOW DO WE FIND THE LENGTH OF THE STRING? WE WILL EXAMINE VARIOUS METHODS, ALGORITHMS, AND CONSIDERATIONS INVOLVED IN DETERMINING THE LENGTH EFFICIENTLY AND ACCURATELY.

UNDERSTANDING THE PROBLEM STATEMENT

WHAT IS THE INPUT?

- A STRING S COMPOSED OF LOWERCASE LETTERS FROM 'A' TO 'Z'.
- THE LENGTH OF THE STRING, N, WHICH INDICATES THE NUMBER OF CHARACTERS IN S.

WHAT IS THE TASK?

- TO COMPUTE AND RETURN THE LENGTH OF THE STRING S.

WHY IS THIS IMPORTANT?

- STRING LENGTH CALCULATION IS FUNDAMENTAL FOR:
- VALIDATING INPUT CONSTRAINTS.
- LOOPING THROUGH ALL CHARACTERS.
- PERFORMING SUBSTRING OPERATIONS.
- IMPLEMENTING ALGORITHMS THAT DEPEND ON STRING SIZE.

METHODS TO FIND THE LENGTH OF A STRING

THERE ARE MULTIPLE APPROACHES TO DETERMINE THE LENGTH OF A STRING, DEPENDING ON THE PROGRAMMING LANGUAGE AND CONTEXT. BELOW, WE EXPLORE THE MOST COMMON METHODS.

1. BUILT-IN FUNCTIONS

MOST PROGRAMMING LANGUAGES PROVIDE A STANDARD FUNCTION OR METHOD TO GET THE LENGTH OF A STRING.

- PYTHON: `len(S)`
- JAVA: `S.length()`
- C++: `S.size()` OR `S.length()`
- JAVASCRIPT: `S.length`

ADVANTAGES:

- VERY EFFICIENT.
- SIMPLE TO IMPLEMENT.
- OPTIMIZED UNDER-THE-HOOD.

EXAMPLE IN PYTHON:

```
"""PYTHON
S = "HELLO"
LENGTH = len(S)
PRINT(LENGTH) OUTPUT: 5
"""
```

2. MANUAL COUNTING

IN SCENARIOS WHERE BUILT-IN FUNCTIONS ARE UNAVAILABLE, OR FOR EDUCATIONAL PURPOSES, YOU CAN MANUALLY COUNT THE CHARACTERS.

ALGORITHM STEPS:

1. INITIALIZE A COUNTER VARIABLE TO ZERO.
2. LOOP THROUGH EACH CHARACTER IN THE STRING.
3. INCREMENT THE COUNTER FOR EACH CHARACTER ENCOUNTERED.
4. RETURN THE COUNTER AS THE LENGTH.

PSEUDOCODE:

```
"""
FUNCTION FINDLENGTH(S):
    COUNT = 0
    FOR EACH CHARACTER IN S:
        COUNT = COUNT + 1
    RETURN COUNT
"""
```

IMPLEMENTATION IN PYTHON:

```
"""PYTHON
DEF FIND_LENGTH(S):
    COUNT = 0
    FOR CH IN S:
        COUNT += 1
    RETURN COUNT

S = "WORLD"
PRINT(FIND_LENGTH(S)) OUTPUT: 5
"""
```

ADVANTAGES:

- DOESN'T RELY ON LANGUAGE-SPECIFIC FUNCTIONS.
- GOOD FOR UNDERSTANDING BASIC STRING TRAVERSAL.

3. RECURSIVE APPROACH

ANOTHER METHOD INVOLVES RECURSION, WHERE THE LENGTH IS DETERMINED BY REDUCING THE STRING STEP-BY-STEP.

ALGORITHM:

- IF THE STRING IS EMPTY, RETURN 0.
- OTHERWISE, RETURN 1 PLUS THE LENGTH OF THE SUBSTRING EXCLUDING THE FIRST CHARACTER.

PSEUDOCODE:

```

FUNCTION RECURSIVELength(S):

IF S IS EMPTY:

RETURN 0

ELSE:

RETURN 1 + RECURSIVELength(S[1:])

```

NOTE: WHILE EDUCATIONAL, RECURSION ISN'T EFFICIENT FOR LARGE STRINGS DUE TO CALL STACK LIMITATIONS.

COMPLEXITY ANALYSIS

UNDERSTANDING THE TIME AND SPACE COMPLEXITY OF STRING LENGTH OPERATIONS HELPS OPTIMIZE PERFORMANCE.

BUILT-IN FUNCTIONS

- TIME COMPLEXITY: $O(1)$ IN MOST LANGUAGES, AS LENGTH IS STORED INTERNALLY.
- SPACE COMPLEXITY: $O(1)$.

MANUAL COUNTING

- TIME COMPLEXITY: $O(N)$, WHERE N IS THE LENGTH OF THE STRING.
- SPACE COMPLEXITY: $O(1)$.

RECURSIVE APPROACH

- TIME COMPLEXITY: $O(N)$, AS IT TRAVERSES ALL CHARACTERS.
- SPACE COMPLEXITY: $O(N)$ DUE TO RECURSION CALL STACK.

PRACTICAL APPLICATIONS OF STRING LENGTH CALCULATION

KNOWING THE LENGTH OF A STRING IS ESSENTIAL ACROSS VARIOUS PROGRAMMING TASKS AND ALGORITHMS.

1. VALIDATING INPUT DATA

- ENSURING STRINGS MEET MINIMUM OR MAXIMUM LENGTH CONSTRAINTS.
- EXAMPLE: PASSWORD VALIDATION REQUIRING A MINIMUM LENGTH.

2. LOOPING AND ITERATION

- TRAVERSING EACH CHARACTER FOR OPERATIONS LIKE ENCRYPTION, PATTERN MATCHING, OR DATA EXTRACTION.

3. SUBSTRING OPERATIONS

- EXTRACTING PARTS OF STRINGS BASED ON THEIR LENGTH.
- EXAMPLE: `'S[0:N]'` TO GET THE ENTIRE STRING.

4. STRING COMPARISON AND SORTING

- COMPARING LENGTHS TO SORT STRINGS BY SIZE.

5. ALGORITHMIC CHALLENGES

- PROBLEMS LIKE FINDING THE LONGEST SUBSTRING WITHOUT REPEATING CHARACTERS DEPEND ON STRING LENGTH.

EDGE CASES AND CONSIDERATIONS

WHEN IMPLEMENTING STRING LENGTH CALCULATION, CONSIDER THE FOLLOWING SCENARIOS.

EMPTY STRING

- STRING `S = ""` (LENGTH = 0). ENSURE YOUR METHOD HANDLES THIS CASE CORRECTLY.

NON-STANDARD CHARACTERS

- ALTHOUGH THE PROBLEM SPECIFIES LOWERCASE LETTERS, REAL-WORLD DATA MAY CONTAIN OTHER CHARACTERS; THE METHOD REMAINS THE SAME.

LARGE STRINGS

- EFFICIENCY BECOMES CRITICAL WITH VERY LARGE N.
- BUILT-IN FUNCTIONS ARE OPTIMIZED FOR SUCH CASES.

IMMUTABLE STRINGS

- IN LANGUAGES LIKE PYTHON, STRINGS ARE IMMUTABLE, SO OPERATIONS LIKE SLICING CREATE NEW STRINGS, IMPACTING PERFORMANCE.

IMPLEMENTING STRING LENGTH CALCULATION IN DIFFERENT LANGUAGES

TO GIVE A CLEARER PERSPECTIVE, HERE ARE CODE SNIPPETS IN VARIOUS POPULAR PROGRAMMING LANGUAGES.

PYTHON

```
'''PYTHON
DEF GET_LENGTH(S):
RETURN LEN(S)
'''
```

JAVA

```
'''JAVA
PUBLIC INT GETSTRINGLENGTH(STRING S) {
RETURN S.LENGTH();
}
'''
```

C++

```
'''CPP
INCLUDE
INT GETLENGTH(CONST STD::STRING& S) {
RETURN S.SIZE();
}
'''
```

JAVASCRIPT

```
'''JAVASCRIPT
FUNCTION GETLENGTH(S) {
RETURN S.LENGTH;
}
'''
```

CONCLUSION

DETERMINING THE LENGTH OF A STRING IS ONE OF THE MOST FUNDAMENTAL OPERATIONS IN PROGRAMMING. WHETHER USING BUILT-IN FUNCTIONS FOR EFFICIENCY AND SIMPLICITY OR IMPLEMENTING MANUAL OR RECURSIVE METHODS FOR EDUCATIONAL PURPOSES, UNDERSTANDING HOW TO FIND THE LENGTH OF A STRING IS ESSENTIAL. IT FORMS THE BASIS FOR MORE COMPLEX

STRING MANIPULATIONS, VALIDATIONS, AND ALGORITHMIC SOLUTIONS. ALWAYS CONSIDER THE CONTEXT, LANGUAGE FEATURES, AND PERFORMANCE IMPLICATIONS WHEN CHOOSING YOUR APPROACH.

BY MASTERING STRING LENGTH CALCULATION, YOU LAY THE GROUNDWORK FOR EFFECTIVE TEXT PROCESSING AND ALGORITHM DEVELOPMENT, WHICH ARE INTEGRAL TO A WIDE ARRAY OF SOFTWARE APPLICATIONS.

FREQUENTLY ASKED QUESTIONS

WHAT IS THE PROBLEM ASKING WHEN GIVEN A STRING S OF LOWERCASE LETTERS?

THE PROBLEM IS ASKING TO DETERMINE THE LENGTH OF A SPECIFIC PROPERTY OR PATTERN WITHIN THE STRING, SUCH AS THE LENGTH OF THE LONGEST SUBSTRING, THE LENGTH OF THE LONGEST SUBSTRING WITHOUT REPEATING CHARACTERS, OR SIMILAR METRICS BASED ON THE CONTEXT.

HOW CAN I FIND THE LENGTH OF THE LONGEST SUBSTRING WITHOUT REPEATING CHARACTERS IN STRING S ?

YOU CAN USE A SLIDING WINDOW APPROACH WITH A HASH SET OR MAP TO KEEP TRACK OF CHARACTERS IN THE CURRENT WINDOW, EXPANDING AND SHRINKING THE WINDOW TO ENSURE ALL CHARACTERS ARE UNIQUE, AND RECORD THE MAXIMUM LENGTH ENCOUNTERED.

WHAT IS THE TIME COMPLEXITY OF FINDING THE LENGTH OF THE LONGEST REPEATED SUBSTRING IN S ?

USING EFFICIENT ALGORITHMS LIKE SUFFIX ARRAYS OR SUFFIX TREES, THE TIME COMPLEXITY CAN BE REDUCED TO $O(N \log N)$ OR BETTER, BUT NAIVE APPROACHES MAY TAKE $O(N^2)$.

CAN DYNAMIC PROGRAMMING BE USED TO SOLVE THIS PROBLEM?

YES, DYNAMIC PROGRAMMING TECHNIQUES CAN BE APPLIED, ESPECIALLY FOR PROBLEMS LIKE FINDING THE LONGEST PALINDROMIC SUBSTRING OR OTHER PATTERN-BASED LENGTHS WITHIN S .

WHAT DATA STRUCTURES ARE USEFUL FOR COMPUTING THE LENGTH OF SUBSTRINGS OR PATTERNS IN S ?

HASH MAPS, HASH SETS, SUFFIX TREES, SUFFIX ARRAYS, AND SLIDING WINDOW DATA STRUCTURES ARE COMMONLY USED TO EFFICIENTLY COMPUTE SUBSTRING LENGTHS AND RELATED METRICS.

HOW DO I HANDLE LARGE STRINGS EFFICIENTLY WHEN CALCULATING SUBSTRING LENGTHS?

EMPLOY OPTIMIZED ALGORITHMS SUCH AS SUFFIX TREES, SUFFIX ARRAYS, OR SLIDING WINDOW TECHNIQUES WITH PROPER PRUNING AND DATA STRUCTURES TO REDUCE TIME AND SPACE COMPLEXITY FOR LARGE INPUT STRINGS.

WHAT ARE COMMON PITFALLS TO AVOID WHEN SOLVING SUBSTRING LENGTH PROBLEMS?

COMMON PITFALLS INCLUDE NOT UPDATING THE WINDOW CORRECTLY, FORGETTING TO HANDLE EDGE CASES LIKE EMPTY STRINGS, OR USING INEFFICIENT ALGORITHMS THAT LEAD TO HIGH TIME COMPLEXITY. ENSURING PROPER INITIALIZATION AND BOUNDARY CHECKS IS ALSO CRUCIAL.

ADDITIONAL RESOURCES

YOU ARE GIVEN A STRING S CONSISTING OF N LOWERCASE LETTERS OF THE ENGLISH ALPHABET. FIND THE LENGTH OF

INTRODUCTION

IN THE REALM OF STRING PROCESSING AND ALGORITHM DESIGN, PROBLEMS INVOLVING STRINGS ARE BOTH FUNDAMENTAL AND DIVERSE. THEY RANGE FROM SIMPLE CHARACTER COUNTING TO COMPLEX PATTERN MATCHING AND SUBSTRING ANALYSIS. ONE SUCH INTRIGUING PROBLEM INVOLVES ANALYZING A STRING (S) OF LENGTH (N) COMPOSED OF LOWERCASE ENGLISH LETTERS, WITH THE GOAL OF DETERMINING A PARTICULAR LENGTH BASED ON SPECIFIC CRITERIA.

THE PHRASE “FIND THE LENGTH OF”—THOUGH INCOMPLETE IN ISOLATION—SUGGESTS A CLASS OF PROBLEMS WHERE THE OBJECTIVE IS TO IDENTIFY THE LENGTH OF A CERTAIN SUBSTRING, SUBSEQUENCE, OR PATTERN THAT MEETS PREDEFINED CONDITIONS. THESE PROBLEMS ARE COMMON IN PROGRAMMING COMPETITIONS, DATA ANALYSIS, AND SOFTWARE DEVELOPMENT, AND THEY OFTEN SERVE AS EXCELLENT EXERCISES FOR UNDERSTANDING ALGORITHMS, DATA STRUCTURES, AND OPTIMIZATION TECHNIQUES.

IN THIS COMPREHENSIVE REVIEW, WE WILL EXPLORE THE TYPICAL PROBLEM SPACE WHERE GIVEN A STRING (S) , THE TASK IS TO FIND THE LENGTH OF A SPECIFIC SEGMENT OR STRUCTURE WITHIN (S) . WE WILL ANALYZE VARIOUS PROBLEM TYPES, ALGORITHMS, AND IMPLEMENTATION STRATEGIES, PROVIDING INSIGHTS INTO THEIR COMPLEXITIES AND NUANCES.

UNDERSTANDING THE CORE PROBLEM

THE BASIC PREMISE

GIVEN A STRING (S) OF LENGTH (N) , COMPOSED OF LOWERCASE ALPHABET CHARACTERS ('a' TO 'z'), THE GOAL IS TO DETERMINE THE LENGTH OF A PARTICULAR SUBSTRING OR SUBSEQUENCE THAT SATISFIES CERTAIN CONSTRAINTS. COMMON OBJECTIVES INCLUDE:

- FINDING THE LENGTH OF THE LONGEST SUBSTRING WITH CERTAIN PROPERTIES.
- DETERMINING THE LENGTH OF THE SHORTEST SUBSTRING CONTAINING ALL CHARACTERS OF A CERTAIN SET.
- COMPUTING THE LENGTH OF THE LONGEST PALINDROMIC SUBSTRING.
- FINDING THE LENGTH OF THE LONGEST SUBSTRING WITH UNIQUE CHARACTERS.
- CALCULATING THE LENGTH OF THE LONGEST SUBSTRING OR SUBSEQUENCE THAT ADHERES TO SPECIFIC PATTERN RULES.

POSSIBLE VARIATIONS

THE PROBLEM'S SCOPE IS BROAD, BUT TYPICAL VARIATIONS INCLUDE:

- LONGEST SUBSTRING WITHOUT REPEATING CHARACTERS: FIND THE MAXIMUM LENGTH OF A SUBSTRING WHERE ALL CHARACTERS ARE UNIQUE.
- LONGEST PALINDROMIC SUBSTRING: FIND THE MAXIMUM LENGTH OF A SUBSTRING THAT READS THE SAME BACKWARD AND FORWARD.
- SUBSTRING WITH ALL CHARACTERS: FIND THE SHORTEST SUBSTRING THAT CONTAINS ALL THE CHARACTERS OF THE ALPHABET OR A SPECIFIED SUBSET.
- MAXIMUM LENGTH OF REPEATING CHARACTERS: FIND THE LONGEST SEQUENCE OF CONSECUTIVE IDENTICAL CHARACTERS.
- LONGEST SUBSTRING WITH K DISTINCT CHARACTERS: FIND THE MAXIMUM LENGTH OF A SUBSTRING THAT CONTAINS AT MOST K DISTINCT CHARACTERS.
- LONGEST SUBSTRING WITH NO FORBIDDEN PATTERNS: FOR EXAMPLE, NO TWO CONSECUTIVE VOWELS, OR NO REPEATED SUBSTRINGS OF LENGTH L .

FORMALIZING THE PROBLEM

WITHOUT LOSS OF GENERALITY, CONSIDER THE PROBLEM:

> GIVEN A STRING (S) OF LENGTH (N) , FIND THE LENGTH OF THE LONGEST SUBSTRING THAT SATISFIES A CERTAIN PROPERTY (P) .

THE PROPERTY (P) VARIES DEPENDING ON THE SPECIFIC PROBLEM, BUT THE CORE CHALLENGE REMAINS: EFFICIENTLY IDENTIFYING THE MAXIMUM LENGTH OF A SEGMENT WITHIN (S) THAT MEETS (P) .

FUNDAMENTAL TECHNIQUES AND ALGORITHMS

TO ADDRESS THESE PROBLEMS EFFECTIVELY, DEVELOPERS AND ALGORITHM DESIGNERS RELY ON A SUITE OF STRATEGIES. HERE, WE EXPLORE THE MOST COMMON AND POWERFUL METHODS.

SLIDING WINDOW TECHNIQUE

THE SLIDING WINDOW IS A PIVOTAL APPROACH WHEN DEALING WITH SUBSTRING PROBLEMS, ESPECIALLY WHERE THE PROPERTY (P) INVOLVES CONSTRAINTS ON CHARACTERS OR SUBSTRINGS.

CONCEPT:

- USE TWO POINTERS (OR INDICES), TYPICALLY CALLED LEFT AND RIGHT.
- EXPAND THE RIGHT POINTER TO INCLUDE MORE CHARACTERS.
- SHRINK THE WINDOW FROM THE LEFT TO MAINTAIN THE PROPERTY (P) .
- TRACK THE MAXIMUM WINDOW SIZE SATISFYING (P) .

APPLICATIONS:

- LONGEST SUBSTRING WITHOUT REPEATING CHARACTERS.
- LONGEST SUBSTRING WITH AT MOST K DISTINCT CHARACTERS.
- LONGEST SUBSTRING WITH A CERTAIN CHARACTER FREQUENCY.

EXAMPLE:

LONGEST SUBSTRING WITHOUT REPEATING CHARACTERS

- INITIALIZE BOTH POINTERS AT 0.
- USE A SET OR MAP TO TRACK CHARACTERS IN THE CURRENT WINDOW.
- EXPAND RIGHT POINTER; IF A DUPLICATE IS FOUND, MOVE THE LEFT POINTER UNTIL THE DUPLICATE IS REMOVED.
- KEEP UPDATING THE MAXIMUM LENGTH FOUND.

TWO-POINTER APPROACH

CLOSELY RELATED TO SLIDING WINDOW, THE TWO-POINTER METHOD IS VERSATILE FOR LINEAR SCANS OF A STRING, ESPECIALLY WHEN THE PROBLEM INVOLVES PAIRWISE RELATIONSHIPS OR SYMMETRY.

PREFIX AND SUFFIX ARRAYS

USEFUL FOR PROBLEMS REQUIRING QUICK RANGE QUERIES, SUCH AS:

- LONGEST PALINDROMIC SUBSTRING: MANACHER'S ALGORITHM.
- LONGEST COMMON SUBSTRING: SUFFIX TREES OR SUFFIX ARRAYS.
- CHARACTER FREQUENCY ANALYSIS: PREFIX SUMS OR PREFIX FREQUENCY ARRAYS.

DYNAMIC PROGRAMMING (DP)

DP IS OFTEN EMPLOYED WHEN THE PROBLEM INVOLVES OVERLAPPING SUBPROBLEMS, SUCH AS:

- LONGEST PALINDROMIC SUBSTRING.
- LONGEST COMMON SUBSEQUENCE.
- LONGEST REPEATING SUBSTRING WITH SPECIFIC CONSTRAINTS.

EXAMPLE:

LONGEST PALINDROMIC SUBSTRING

- USE A DP TABLE WHERE $dp[i][j]$ INDICATES WHETHER $S[i..j]$ IS A PALINDROME.
- FILL THE TABLE ITERATIVELY BASED ON SMALLER SUBSTRINGS.

HASHING AND STRING MATCHING ALGORITHMS

STRING HASHING (E.G., RABIN-KARP) ALLOWS EFFICIENT SUBSTRING COMPARISONS, ESPECIALLY USEFUL FOR:

- DETECTING REPEATED SUBSTRINGS.
- PATTERN MATCHING.
- IDENTIFYING DUPLICATE SUBSTRINGS EFFICIENTLY.

SUFFIX TREES AND SUFFIX ARRAYS

ADVANCED DATA STRUCTURES TO SOLVE COMPLEX PROBLEMS LIKE:

- LONGEST REPEATED SUBSTRING.
- LONGEST SUBSTRING WITH CERTAIN SUBSTRING PROPERTIES.
- STRING COMPRESSION AND PATTERN DETECTION.

DEEP DIVE INTO KEY PROBLEMS AND SOLUTIONS

LONGEST SUBSTRING WITHOUT REPEATING CHARACTERS

PROBLEM STATEMENT:

GIVEN A STRING S , FIND THE LENGTH OF THE LONGEST SUBSTRING WITHOUT CHARACTER REPETITION.

APPROACH:

- USE A SLIDING WINDOW.
- MAINTAIN A SET OR HASH MAP TO KEEP TRACK OF CHARACTERS IN THE CURRENT WINDOW.
- EXPAND THE WINDOW BY MOVING THE RIGHT POINTER.
- IF A DUPLICATE IS ENCOUNTERED, MOVE THE LEFT POINTER UNTIL THE DUPLICATE IS REMOVED.
- TRACK AND UPDATE THE MAXIMUM WINDOW SIZE THROUGHOUT.

IMPLEMENTATION:

```

"""PYTHON
def lengthOfLongestSubstring(s):
    seen = set()
    left = 0
    max_length = 0
    for right in range(len(s)):
        while s[right] in seen:
            seen.remove(s[left])
            left += 1
        seen.add(s[right])
        max_length = max(max_length, right - left + 1)
    return max_length
"""

```

TIME COMPLEXITY:

- $O(N)$, SINCE EACH CHARACTER IS VISITED AT MOST TWICE.

LONGEST PALINDROMIC SUBSTRING

PROBLEM:

FIND THE LONGEST SUBSTRING WITHIN S THAT READS THE SAME BACKWARD AND FORWARD.

APPROACHES:

1. EXPAND AROUND CENTER:

- FOR EACH POSITION, EXPAND OUTWARD TO FIND PALINDROMES OF ODD AND EVEN LENGTH.
- TRACK THE MAXIMUM PALINDROME LENGTH FOUND.

2. DYNAMIC PROGRAMMING:

- USE A TABLE TO MARK WHETHER SUBSTRINGS ARE PALINDROMES.
- FILL THE TABLE ITERATIVELY, STARTING FROM SHORTER SUBSTRINGS.

3. MANACHER'S ALGORITHM:

- A LINEAR-TIME ALGORITHM FOR FINDING THE LONGEST PALINDROMIC SUBSTRING.

IMPLEMENTATION (EXPAND AROUND CENTER):

```
"""PYTHON
DEF LONGESTPALINDROME(S):
    MAX_LEN = 1
    START = 0
    FOR I IN RANGE(LEN(S)):
        ODD LENGTH PALINDROME
        L, R = I, I
        WHILE L >= 0 AND R < LEN(S) AND S[L] == S[R]:
            IF R - L + 1 > MAX_LEN:
                MAX_LEN = R - L + 1
                START = L
            L -= 1
            R += 1
        EVEN LENGTH PALINDROME
        L, R = I, I + 1
        WHILE L >= 0 AND R < LEN(S) AND S[L] == S[R]:
            IF R - L + 1 > MAX_LEN:
                MAX_LEN = R - L + 1
                START = L
            L -= 1
            R += 1
    RETURN S[START:START+MAX_LEN]
"""
```

COMPLEXITY:

- $(O(N^2))$, MANAGEABLE FOR MODERATE (N) .

LONGEST SUBSTRING WITH K DISTINCT CHARACTERS

PROBLEM:

GIVEN (S) AND AN INTEGER (K) , FIND THE MAXIMUM LENGTH OF A SUBSTRING THAT CONTAINS AT MOST (K) DISTINCT CHARACTERS.

APPROACH:

- SLIDING WINDOW WITH A FREQUENCY MAP.
- EXPAND THE RIGHT POINTER, UPDATING THE COUNT OF CHARACTERS.
- SHRINK FROM THE LEFT WHEN THE NUMBER OF DISTINCT CHARACTERS EXCEEDS (K) .
- TRACK MAXIMUM LENGTH DURING THE PROCESS.

IMPLEMENTATION:

```
"""PYTHON
DEF LENGTHOFLONGESTSUBSTRINGKDISTINCT(S, K):
    FROM COLLECTIONS IMPORT DEFAULTDICT
    FREQ = DEFAULTDICT(INT)
    LEFT = 0
    MAX_LENGTH = 0
    DISTINCT_COUNT = 0
```

```

FOR RIGHT IN RANGE(LEN(S)):
    IF FREQ[S[RIGHT]] == 0:
        DISTINCT_COUNT += 1
        FREQ[S[RIGHT]] += 1

    WHILE DISTINCT_COUNT > K:
        FREQ[S[LEFT]] -= 1
        IF FREQ[S[LEFT]] == 0:
            DISTINCT_COUNT -= 1
            LEFT += 1

    MAX_LENGTH = MAX(MAX_LENGTH, RIGHT - LEFT + 1)
RETURN MAX_LENGTH
'''

```

TIME COMPLEXITY:
 - $O(N)$, SINCE EACH CHARACTER IS PROCESSED AT MOST TWICE.

ADVANCED TOPICS AND DATA STRUCTURES

SUFFIX TREES AND SUFFIX ARRAYS

FOR MORE COMPLEX SUBSTRING PROBLEMS, SUCH AS:

- LONGEST REPEATED SUBSTRING.
- LONGEST SUBSTRING APPEARING MORE THAN ONCE.
- PATTERN MATCHING WITH MULTIPLE PATTERNS.

SUFFIX TREES AND SUFFIX ARRAYS PROVIDE EFFICIENT SOLUTIONS.

- SUFFIX TREE: A COMPRESSED TRIE OF ALL

Ou Are Given A String S Consisting Of N Lowercase Letters Of The English Alphabet Find The Length Of

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