

CREATE AN BST TREE BY ADDING THE FOLLOWING NUMBERS TO THESE TREES IN ORDER GIVEN: 80,70,60,50,40,30,20,10,

CREATE AN BST TREE BY ADDING THE FOLLOWING NUMBERS TO THESE TREES IN ORDER GIVEN:
80,70,60,50,40,30,20,10,

INTRODUCTION TO BINARY SEARCH TREES (BSTs)

WHAT IS A BINARY SEARCH TREE?

A BINARY SEARCH TREE (BST) IS A SPECIALIZED TYPE OF BINARY TREE WHERE EACH NODE FOLLOWS A SPECIFIC ORDER:

- THE LEFT SUBTREE OF A NODE CONTAINS ONLY NODES WITH VALUES LESS THAN THE NODE'S VALUE.
- THE RIGHT SUBTREE CONTAINS ONLY NODES WITH VALUES GREATER THAN THE NODE'S VALUE.

THIS PROPERTY FACILITATES EFFICIENT SEARCHING, INSERTION, AND DELETION OPERATIONS, OFTEN OPERATING IN LOGARITHMIC TIME COMPLEXITY FOR BALANCED TREES.

IMPORTANCE OF BSTs IN COMPUTER SCIENCE

BSTs SERVE AS FUNDAMENTAL DATA STRUCTURES IN VARIOUS APPLICATIONS, INCLUDING:

- DATABASE INDEXING
- IMPLEMENTING ASSOCIATIVE ARRAYS (LIKE MAPS AND SETS)
- SORTING ALGORITHMS
- PRIORITY QUEUE IMPLEMENTATIONS

UNDERSTANDING HOW BSTs ARE CONSTRUCTED, ESPECIALLY HOW INSERTION ORDER INFLUENCES THEIR SHAPE, IS CRUCIAL FOR OPTIMIZING THEIR PERFORMANCE.

STEP-BY-STEP CONSTRUCTION OF THE BST WITH GIVEN SEQUENCE

SEQUENCE OF NUMBERS TO INSERT

THE NUMBERS ARE INSERTED INTO THE BST IN THE FOLLOWING ORDER:

- 80, 70, 60, 50, 40, 30, 20, 10

INITIAL INSERTION: 80

- SINCE THE TREE IS EMPTY, 80 BECOMES THE ROOT NODE.
- TREE NOW:

'''

80

'''

INSERT 70

- 70 IS LESS THAN 80, SO IT BECOMES THE LEFT CHILD OF 80.
- TREE NOW:

'''

```
80
/
70
'''
```

INSERT 60

- $60 < 80$? GO LEFT
- $60 < 70$? GO LEFT OF 70
- INSERTED AS LEFT CHILD OF 70
- TREE:

'''

```
80
/
70
/
60
'''
```

INSERT 50

- $50 < 80$? LEFT
- $50 < 70$? LEFT
- $50 < 60$? LEFT
- INSERTED AS LEFT CHILD OF 60
- TREE:

'''

```
80
/
70
/
60
/
50
'''
```

INSERT 40

- $40 < 80$? LEFT
- $40 < 70$? LEFT
- $40 < 60$? LEFT
- $40 < 50$? LEFT
- INSERTED AS LEFT CHILD OF 50
- TREE:

""
80
/
70
/
60
/
50
/
40
""

INSERT 30

- $30 < 80$? LEFT
- $30 < 70$? LEFT
- $30 < 60$? LEFT
- $30 < 50$? LEFT
- $30 < 40$? LEFT
- INSERTED AS LEFT CHILD OF 40
- TREE:

""
80
/
70
/
60
/
50
/
40
/
30
""

INSERT 20

- $20 < 80$? LEFT
- $20 < 70$? LEFT
- $20 < 60$? LEFT
- $20 < 50$? LEFT
- $20 < 40$? LEFT
- $20 < 30$? LEFT
- INSERTED AS LEFT CHILD OF 30
- TREE:

""
80
/
70
/
60
/
50
/
40
""

/
30
/
20
'''

INSERT 10

- $10 < 80$? LEFT
- $10 < 70$? LEFT
- $10 < 60$? LEFT
- $10 < 50$? LEFT
- $10 < 40$? LEFT
- $10 < 30$? LEFT
- $10 < 20$? LEFT
- INSERTED AS LEFT CHILD OF 20
- FINAL TREE:

'''
80
/
70
/
60
/
50
/
40
/
30
/
20
/
10
'''

UNDERSTANDING THE SHAPE OF THE TREE

TREE CHARACTERISTICS

- THE TREE IS SKEWED TO THE LEFT, RESEMBLING A LINKED LIST.
- EACH INSERTED NUMBER IS LESS THAN THE PREVIOUS, LEADING TO A DEGENERATE STRUCTURE.
- THE DEPTH OF THE TREE IS EQUAL TO THE NUMBER OF INSERTED ELEMENTS MINUS ONE.

IMPACT OF INSERTION ORDER

- INSERTING DECREASING NUMBERS CAUSES THE TREE TO BECOME SKEWED.
- THIS REDUCES THE EFFICIENCY OF SEARCH OPERATIONS, WHICH DEGRADE FROM $O(\log N)$ TO $O(N)$ IN THE WORST CASE.
- TO MAINTAIN BALANCE, DIFFERENT INSERTION STRATEGIES OR SELF-BALANCING TREES LIKE AVL OR RED-BLACK TREES ARE USED.

VISUAL REPRESENTATION OF FINAL TREE

```
""
80
/
70
/
60
/
50
/
40
/
30
/
20
/
10
""
```

OPERATIONS ON THE CONSTRUCTED BST

SEARCHING FOR ELEMENTS

- TO FIND A NUMBER, START AT THE ROOT AND TRAVERSE LEFT OR RIGHT DEPENDING ON VALUE COMPARISONS.
- DUE TO THE SKEWED STRUCTURE, SEARCH OPERATIONS MAY REQUIRE TRAVERSING ALMOST ALL NODES.

INSERTING ADDITIONAL ELEMENTS

- INSERTING LARGER NUMBERS WOULD NORMALLY BALANCE THE TREE BETTER.
- SINCE THE CURRENT TREE IS SKEWED, INSERTING LARGER NUMBERS WOULD GO TO THE RIGHT, MAINTAINING OR WORSENING IMBALANCE.

DELETING NODES

- DELETION INVOLVES SEVERAL CASES:
 1. NODE WITH NO CHILDREN (LEAF NODE): SIMPLY REMOVE.
 2. NODE WITH ONE CHILD: REPLACE NODE WITH ITS CHILD.
 3. NODE WITH TWO CHILDREN: REPLACE WITH IN-ORDER SUCCESSOR OR PREDECESSOR.
- DUE TO THE SKEWED STRUCTURE, DELETIONS AND REBALANCING MIGHT BE MORE COMPLEX.

STRATEGIES TO IMPROVE TREE BALANCE

SELF-BALANCING BSTs

- TO PREVENT SKEWED TREES, SELF-BALANCING ALGORITHMS ARE EMPLOYED:

- AVL TREES: MAINTAIN BALANCE FACTORS AND ROTATE NODES AS NECESSARY.
- RED-BLACK TREES: USE COLOR PROPERTIES AND ROTATIONS TO ENSURE APPROXIMATE BALANCE.
- THESE STRUCTURES ENSURE $O(\log N)$ HEIGHT REGARDLESS OF INSERTION ORDER.

ORDER OF INSERTION MATTERS

- INSERTING NUMBERS IN A SPECIFIC ORDER CAN LEAD TO MORE BALANCED TREES.
- FOR EXAMPLE, INSERTING THE MIDDLE ELEMENT FIRST, THEN RECURSIVELY INSERTING MIDPOINTS OF SUBARRAYS, PRODUCES A BALANCED BST.

CONSTRUCTING A BALANCED BST FROM SORTED ARRAY

- IF THE DATA IS SORTED, A BALANCED BST CAN BE BUILT EFFICIENTLY:
 1. SELECT THE MIDDLE ELEMENT AS ROOT.
 2. RECURSIVELY BUILD LEFT AND RIGHT SUBTREES FROM SUBARRAYS.
- THIS STRATEGY MINIMIZES TREE HEIGHT AND OPTIMIZES SEARCH OPERATIONS.

CONCLUSION

SUMMARY OF THE CONSTRUCTION PROCESS

- THE SEQUENCE OF INSERTING 80,70,60,50,40,30,20,10 RESULTS IN A HIGHLY SKEWED, LEFT-DEGENERATE BST.
- WHILE STRAIGHTFORWARD TO IMPLEMENT, SUCH A STRUCTURE IS INEFFICIENT FOR SEARCH AND OTHER OPERATIONS.

LESSONS LEARNED

- THE ORDER OF INSERTION CRITICALLY INFLUENCES THE SHAPE AND PERFORMANCE OF A BST.
- FOR OPTIMAL PERFORMANCE, CONSIDER BALANCING TECHNIQUES OR DIFFERENT INSERTION STRATEGIES.
- UNDERSTANDING THE STRUCTURE AND BEHAVIOR OF BSTS HELPS IN DESIGNING EFFICIENT ALGORITHMS AND DATA STRUCTURES.

FINAL THOUGHTS

CONSTRUCTING A BST WITH A SPECIFIC SEQUENCE OFFERS INSIGHT INTO HOW INSERTION ORDER IMPACTS ITS SHAPE AND EFFICIENCY. RECOGNIZING THE PITFALLS OF SKEWED TREES UNDERSCORES THE IMPORTANCE OF BALANCING TECHNIQUES IN REAL-WORLD APPLICATIONS. WHETHER BUILDING A SIMPLE BST OR EMPLOYING SELF-BALANCING VARIANTS, UNDERSTANDING THESE FOUNDATIONAL PRINCIPLES IS KEY TO EFFECTIVE SOFTWARE DESIGN AND DATA MANAGEMENT.

FREQUENTLY ASKED QUESTIONS

WHAT IS THE SEQUENCE OF INSERTIONS USED TO CREATE THE BST IN THIS EXAMPLE?

THE NUMBERS ARE INSERTED IN THE FOLLOWING ORDER: 80, 70, 60, 50, 40, 30, 20, 10.

WHAT WILL THE BST LOOK LIKE AFTER INSERTING THE NUMBERS IN THE GIVEN ORDER?

THE BST WILL BE SKEWED TO THE LEFT, WITH EACH NEW NUMBER BECOMING THE LEFT CHILD OF THE PREVIOUS NODE, FORMING A

DESCENDING CHAIN FROM 80 DOWN TO 10.

IS THE RESULTING BST BALANCED AFTER INSERTING THESE NUMBERS IN ORDER?

NO, THE TREE WILL BE HIGHLY UNBALANCED, RESEMBLING A LINKED LIST LEANING TO THE LEFT DUE TO DECREASING INSERTIONS.

HOW DOES INSERTING DECREASING NUMBERS AFFECT THE SHAPE OF A BST?

INSERTING DECREASING NUMBERS TENDS TO PRODUCE A SKEWED, LEFT-LEANING TREE, WHICH CAN DEGRADE SEARCH EFFICIENCY TO LINEAR TIME.

WHAT IS THE INORDER TRAVERSAL OF THE BST AFTER INSERTING THESE NUMBERS?

THE INORDER TRAVERSAL WILL LIST THE NUMBERS IN ASCENDING ORDER: 10, 20, 30, 40, 50, 60, 70, 80.

HOW CAN WE BALANCE THIS SKEWED BST AFTER INSERTING THESE NUMBERS?

TO BALANCE THE TREE, WE CAN PERFORM OPERATIONS LIKE AVL ROTATIONS, OR BUILD A BALANCED TREE FROM SORTED DATA USING TECHNIQUES SUCH AS AVL OR RED-BLACK TREES.

WHAT IS THE TIME COMPLEXITY OF SEARCHING FOR AN ELEMENT IN THIS SKEWED BST?

THE SEARCH TIME COMPLEXITY IN A SKEWED BST IS $O(n)$, WHERE n IS THE NUMBER OF NODES, BECAUSE IT BEHAVES LIKE A LINKED LIST.

CAN INSERTING THESE NUMBERS IN A DIFFERENT ORDER PRODUCE A MORE BALANCED BST?

YES, INSERTING NUMBERS IN A MEDIAN-FIRST OR SHUFFLED ORDER CAN PRODUCE A MORE BALANCED BST, IMPROVING SEARCH PERFORMANCE.

WHY IS IT IMPORTANT TO CONSIDER BALANCING WHEN CREATING BSTS WITH SEQUENTIAL DATA?

BALANCING PREVENTS THE TREE FROM BECOMING SKEWED, MAINTAINING EFFICIENT SEARCH, INSERTION, AND DELETION OPERATIONS WITH AVERAGE $O(\log n)$ COMPLEXITY.

WHAT IS A PRACTICAL APPLICATION OF UNDERSTANDING HOW INSERTION ORDER AFFECTS BST SHAPE?

KNOWING THIS HELPS IN DESIGNING EFFICIENT DATA STRUCTURES, OPTIMIZING SEARCH ALGORITHMS, AND AVOIDING WORST-CASE SCENARIOS LIKE DEGENERATE TREES IN REAL-WORLD APPLICATIONS.

ADDITIONAL RESOURCES

CREATING A BINARY SEARCH TREE (BST) BY SEQUENTIALLY ADDING THE NUMBERS 80, 70, 60, 50, 40, 30, 20, 10

BUILDING A BINARY SEARCH TREE (BST) IS A FUNDAMENTAL TASK IN COMPUTER SCIENCE THAT ENCAPSULATES PRINCIPLES OF DATA ORGANIZATION, EFFICIENCY IN SEARCH OPERATIONS, AND THE ELEGANCE OF RECURSIVE STRUCTURES. IN THIS ARTICLE, WE EXPLORE THE STEP-BY-STEP PROCESS OF CONSTRUCTING A BST BY INSERTING THE SEQUENCE OF NUMBERS: 80, 70, 60, 50, 40, 30, 20, 10, IN THAT SPECIFIC ORDER. WE WILL ANALYZE THE EVOLUTION OF THE TREE WITH EACH INSERTION, EXAMINE ITS STRUCTURAL PROPERTIES, AND DISCUSS THE IMPLICATIONS OF INSERTION ORDER ON TREE SHAPE AND PERFORMANCE.

UNDERSTANDING BINARY SEARCH TREES (BSTs)

WHAT IS A BINARY SEARCH TREE?

A BINARY SEARCH TREE IS A SPECIALIZED BINARY TREE WHERE EACH NODE ADHERES TO THE BINARY SEARCH PROPERTY: THE VALUE STORED IN THE LEFT SUBTREE OF A NODE IS LESS THAN THE NODE'S VALUE, AND THE VALUE IN THE RIGHT SUBTREE IS GREATER THAN THE NODE'S VALUE. THIS PROPERTY ENABLES EFFICIENT SEARCHING, INSERTION, AND DELETION OPERATIONS, OFTEN WITH AVERAGE TIME COMPLEXITIES OF $O(\log n)$, ASSUMING THE TREE IS BALANCED.

CORE OPERATIONS AND THEIR SIGNIFICANCE

- INSERTION: ADDS A NEW NODE WHILE MAINTAINING THE BST PROPERTY.
- SEARCH: FINDS WHETHER A VALUE EXISTS IN THE TREE.
- DELETION: REMOVES A NODE, POTENTIALLY RESTRUCTURING THE TREE TO PRESERVE PROPERTIES.
- TRAVERSAL: VISITS NODES IN A SPECIFIC ORDER (IN-ORDER, PRE-ORDER, POST-ORDER) TO RETRIEVE DATA.

THE SHAPE OF THE BST SIGNIFICANTLY AFFECTS PERFORMANCE. A BALANCED TREE OFFERS NEAR-OPTIMAL OPERATION TIMES, WHILE AN UNBALANCED, SKEWED TREE CAN DEGRADE TO LINEAR TIME.

STEP-BY-STEP CONSTRUCTION OF THE BST WITH THE GIVEN SEQUENCE

TO UNDERSTAND HOW THE TREE EVOLVES, WE EXAMINE EACH INSERTION SEQUENTIALLY, NOTING THE STRUCTURE AT EACH STEP.

INITIAL INSERTION: 80

- THE TREE IS EMPTY INITIALLY.
- INSERT 80 AS THE ROOT NODE.

TREE AFTER INSERTION:

```
""
80
""
```

THIS STARTING POINT IS STRAIGHTFORWARD; THE ROOT IS 80.

SECOND INSERTION: 70

- COMPARE 70 WITH 80.
- SINCE $70 < 80$, INSERT 70 AS THE LEFT CHILD OF 80.

TREE:

```
""
80
/
70
""
```

THE TREE NOW HAS A LEFT SUBTREE UNDER THE ROOT.

THIRD INSERTION: 60

- COMPARE 60 WITH 80: $60 < 80$? GO LEFT.
- COMPARE 60 WITH 70: $60 < 70$? GO LEFT.
- LEFT OF 70 IS EMPTY, INSERT 60 HERE.

TREE:

```
""
80
/
70
/
60
""
```

THE LEFT SUBTREE DEEPENS, FORMING A SKEWED STRUCTURE ON THE LEFT.

FOURTH INSERTION: 50

- $50 < 80$? GO LEFT.
- $50 < 70$? GO LEFT.
- $50 < 60$? GO LEFT.
- INSERT 50 AS THE LEFT CHILD OF 60.

TREE:

```
""
80
/
70
/
60
/
50
""
```

THE LEFT-SKEWED CHAIN CONTINUES, EMPHASIZING AN UNBALANCED STRUCTURE.

FIFTH INSERTION: 40

- TRAVERSE: $40 < 80$? LEFT.
- $40 < 70$? LEFT.
- $40 < 60$? LEFT.
- $40 < 50$? LEFT.
- INSERT 40 AS THE LEFT CHILD OF 50.

TREE:

```
""
80
/
70
/
60
/
50
/
40
""
```

THE PATTERN PERSISTS, CREATING A HIGHLY SKEWED LEFT CHAIN.

SIXTH INSERTION: 30

- FOLLOW THE SAME COMPARISON: $30 < 80$? LEFT.
- $30 < 70$? LEFT.
- $30 < 60$? LEFT.
- $30 < 50$? LEFT.
- $30 < 40$? LEFT.
- INSERT 30 AS THE LEFT CHILD OF 40.

TREE:

```
""
80
/
70
/
60
/
50
/
40
/
30
""
```

THE TREE'S LEFT SUBTREE BECOMES A LINEAR CHAIN OF DECREASING VALUES.

SEVENTH INSERTION: 20

- CONTINUE COMPARISON: $20 < 80$? LEFT.
- $20 < 70$? LEFT.
- $20 < 60$? LEFT.
- $20 < 50$? LEFT.
- $20 < 40$? LEFT.
- $20 < 30$? LEFT.
- INSERT 20 AS THE LEFT CHILD OF 30.

TREE:

```
""  
80  
/  
70  
/  
60  
/  
50  
/  
40  
/  
30  
/  
20  
""
```

THIS PATTERN RESULTS IN A DEEP, LEFT-SKEWED CHAIN.

EIGHT INSERTION: 10

- PROCEED WITH COMPARISONS: $10 < 80, 70, 60, 50, 40, 30, 20$.
- INSERT 10 AS THE LEFT CHILD OF 20.

FINAL TREE STRUCTURE:

```
""  
80  
/  
70  
/  
60  
/  
50  
/  
40  
/  
30  
/  
20  
/  
10  
""
```

THE CUMULATIVE EFFECT OF THESE INSERTIONS IS A HIGHLY UNBALANCED, LINEAR TREE LEANING ENTIRELY TO THE LEFT, RESEMBLING A LINKED LIST.

STRUCTURAL ANALYSIS OF THE CONSTRUCTED BST

SHAPE AND HEIGHT:

- THE RESULTANT TREE IS SKEWED, WITH ALL NODES FORMING A SINGLE LEFTWARD CHAIN.
- HEIGHT OF THE TREE: 8 (NUMBER OF NODES FROM ROOT TO THE DEEPEST LEAF).
- THE NUMBER OF LEVELS EQUALS THE NUMBER OF INSERTED NODES, INDICATING THE WORST-CASE SCENARIO FOR SEARCH OPERATIONS IN THIS PARTICULAR INSTANCE.

IMPLICATIONS OF THE INSERTION ORDER:

- THE SEQUENTIAL INSERTION OF DECREASING NUMBERS INHERENTLY PRODUCES A UNBALANCED, DEGENERATE TREE.
- THIS STRUCTURE NEGATES THE BENEFITS OF BSTs FOR EFFICIENT SEARCH, WHICH IDEALLY OPERATE IN LOGARITHMIC TIME.
- THE SHAPE IS SIMILAR TO A LINKED LIST, LEADING TO $O(n)$ SEARCH TIME IN THE WORST CASE.

CONTRAST WITH BALANCED BSTs:

- INSERTING NUMBERS IN A DIFFERENT ORDER (SUCH AS A RANDOMIZED SEQUENCE OR SORTED IN ASCENDING ORDER) COULD PRODUCE A MORE BALANCED TREE.
- TECHNIQUES LIKE SELF-BALANCING BSTs (E.G., AVL TREES, RED-BLACK TREES) AUTOMATICALLY MAINTAIN BALANCE TO OPTIMIZE PERFORMANCE.

PERFORMANCE CONSIDERATIONS AND PRACTICAL IMPLICATIONS

EFFICIENCY OF OPERATIONS:

- THE SKEWED STRUCTURE SIGNIFICANTLY HAMPERS EFFICIENCY.
- SEARCH, INSERTION, AND DELETION OPERATIONS DEGRADE TO LINEAR TIME COMPLEXITY $O(n)$.

MEMORY AND STORAGE:

- THE LINEAR SHAPE LEADS TO LESS EFFICIENT USE OF MEMORY CACHE, AS TRAVERSAL INVOLVES FOLLOWING A CHAIN OF NODES WITHOUT THE BENEFITS OF LOCALITY OF REFERENCE.

REAL-WORLD APPLICATIONS:

- WHEN CREATING BSTs FROM SEQUENTIAL DATA, INSERTION ORDER IS CRITICAL.
- DATA SHOULD IDEALLY BE INSERTED IN A MANNER THAT PROMOTES BALANCE.
- ALTERNATIVELY, EMPLOYING SELF-BALANCING BSTs CAN MITIGATE ORDER-RELATED ISSUES.

CONCLUSION: LESSONS FROM CONSTRUCTING A LEFT-SKEWED BST

CONSTRUCTING A BINARY SEARCH TREE BY INSERTING A SEQUENCE OF DECREASING NUMBERS—80, 70, 60, 50, 40, 30, 20,

10—SERVES AS A POTENT ILLUSTRATION OF HOW INSERTION ORDER INFLUENCES TREE SHAPE AND PERFORMANCE. THE RESULTING STRUCTURE, A LINEAR CHAIN OF NODES LEANING LEFTWARD, EXEMPLIFIES A WORST-CASE SCENARIO FOR BSTs, UNDERMINING THEIR CORE ADVANTAGE OF LOGARITHMIC SEARCH TIME.

THIS EXERCISE UNDERSCORES THE IMPORTANCE OF MINDFUL DATA INSERTION STRATEGIES AND THE VALUE OF SELF-BALANCING TREES IN REAL-WORLD APPLICATIONS. WHETHER IN DATABASE INDEXING, MEMORY MANAGEMENT, OR ALGORITHM DESIGN, UNDERSTANDING THE NUANCES OF TREE CONSTRUCTION ENABLES DEVELOPERS AND COMPUTER SCIENTISTS TO OPTIMIZE DATA STRUCTURES FOR EFFICIENCY AND ROBUSTNESS. BY ANALYZING THIS SPECIFIC CASE, PRACTITIONERS ARE BETTER EQUIPPED TO RECOGNIZE POTENTIAL PITFALLS AND ADOPT STRATEGIES THAT PRESERVE THE BENEFITS OF BINARY SEARCH TREES IN DIVERSE COMPUTATIONAL CONTEXTS.

Create An Bst Tree By Adding The Following Numbers To These Trees In Order Given 80 70 60 50 40 30 20 10

Create An Bst Tree By Adding The Following Numbers To These Trees In Order Given 80 70 60 50 40 30 20 10

Related Articles

- [Monetary Policy Is Set By TheMultiple Choice A.Regional Federal Reserve Banks. B.Federal Advisory Council.](#)
- [Please Help! A Study Must Be Valid To Be Considered ReliablePlease Select The Best Answer From The Choices](#)
- [Carefully Review All The Pieces Of Evidence For The Theory Of Plate Tectonics, I Would Like An Explanation](#)

[Back to Home](#)