

GIVEN THE MEMORY PARTITIONS OF 600k, 700k, 300k, 200k IN ORDER, HOW WOULD EACH OF THE FIRST-FIT AND BEST-FIT

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WHEN IT COMES TO MEMORY MANAGEMENT IN OPERATING SYSTEMS, THE PROCESS OF ALLOCATING MEMORY EFFICIENTLY IS CRUCIAL FOR OPTIMIZING SYSTEM PERFORMANCE. TWO COMMONLY USED ALGORITHMS FOR MEMORY ALLOCATION ARE FIRST-FIT AND BEST-FIT. BOTH METHODS AIM TO ALLOCATE PROCESSES TO AVAILABLE MEMORY PARTITIONS BUT DIFFER SIGNIFICANTLY IN HOW THEY SELECT SUITABLE PARTITIONS. IN THIS ARTICLE, WE WILL EXPLORE HOW EACH OF THESE ALGORITHMS OPERATES GIVEN A SPECIFIC SET OF MEMORY PARTITIONS: 600k, 700k, 300k, AND 200k, ARRANGED IN THAT ORDER.

OUR GOAL IS TO UNDERSTAND HOW FIRST-FIT AND BEST-FIT ALGORITHMS HANDLE MEMORY ALLOCATION WHEN PRESENTED WITH A SEQUENCE OF PROCESSES REQUIRING DIFFERENT MEMORY SIZES. WE WILL ANALYZE EACH METHOD STEP-BY-STEP, ILLUSTRATING THE ALLOCATION PROCESS AND DISCUSSING THEIR ADVANTAGES AND DISADVANTAGES.

UNDERSTANDING MEMORY ALLOCATION ALGORITHMS

BEFORE DIVING INTO THE SPECIFIC EXAMPLE, LET'S BRIEFLY REVIEW WHAT FIRST-FIT AND BEST-FIT ALGORITHMS ARE, ALONG WITH THEIR CORE PRINCIPLES.

FIRST-FIT ALLOCATION

- PRINCIPLE: ALLOCATE THE FIRST PARTITION ENCOUNTERED IN THE MEMORY LIST THAT IS LARGE ENOUGH TO SATISFY THE PROCESS'S MEMORY REQUEST.
- PROCESS:
 1. START FROM THE BEGINNING OF THE LIST OF PARTITIONS.
 2. ALLOCATE THE PROCESS TO THE FIRST PARTITION THAT IS LARGE ENOUGH.
 3. IF THE PARTITION IS LARGER THAN NEEDED, SPLIT THE PARTITION, AND ALLOCATE THE REQUIRED MEMORY TO THE PROCESS, LEAVING THE REMAINING PART AS A NEW FREE PARTITION.
 4. PROCEED TO THE NEXT PROCESS AND REPEAT THE PROCESS.

ADVANTAGES:

- SIMPLE AND FAST TO IMPLEMENT.
- TENDS TO ALLOCATE MEMORY QUICKLY, REDUCING SEARCH TIME.

DISADVANTAGES:

- CAN LEAD TO FRAGMENTATION OVER TIME.
- MAY LEAVE SMALL UNUSABLE PARTITIONS AT THE BEGINNING OF MEMORY.

BEST-FIT ALLOCATION

- PRINCIPLE: ALLOCATE THE PROCESS TO THE SMALLEST PARTITION THAT IS LARGE ENOUGH TO SATISFY THE REQUEST.
- PROCESS:
 1. SEARCH THE ENTIRE LIST OF FREE PARTITIONS.
 2. FIND THE PARTITION THAT IS CLOSEST IN SIZE TO THE REQUESTED MEMORY BUT STILL LARGE ENOUGH.

3. ALLOCATE THE PROCESS TO THIS PARTITION, SPLITTING IF NECESSARY.
4. CONTINUE FOR SUBSEQUENT PROCESSES.

ADVANTAGES:

- MINIMIZES WASTAGE OF LARGE PARTITIONS.
- REDUCES INTERNAL FRAGMENTATION.

DISADVANTAGES:

- MORE TIME-CONSUMING DUE TO SEARCHING FOR THE BEST FIT.
- CAN LEAD TO EXTERNAL FRAGMENTATION OVER TIME.

MEMORY PARTITION SETUP AND PROCESS REQUESTS

GIVEN THE INITIAL MEMORY PARTITIONS:

- PARTITION 1: 600k
- PARTITION 2: 700k
- PARTITION 3: 300k
- PARTITION 4: 200k

ASSUMING THE PROCESSES ARRIVE WITH THE FOLLOWING MEMORY REQUIREMENTS:

- PROCESS A: 212k
- PROCESS B: 417k
- PROCESS C: 112k
- PROCESS D: 426k

THE ALLOCATION PROCESS WILL PROCEED STEP-BY-STEP FOR EACH ALGORITHM.

MEMORY ALLOCATION USING FIRST-FIT ALGORITHM

LET'S ANALYZE HOW FIRST-FIT WOULD ALLOCATE EACH PROCESS GIVEN THE SEQUENCE OF PARTITIONS.

STEP 1: ALLOCATING PROCESS A (212k)

- START FROM PARTITION 1 (600k): $600k \geq 212k$ ☒ ALLOCATE HERE.
- REMAINING PARTITION SIZE: $600k - 212k = 388k$.
- UPDATED PARTITIONS:
- PARTITION 1: 388k (REMAINING FREE)
- PARTITION 2: 700k
- PARTITION 3: 300k
- PARTITION 4: 200k

STEP 2: ALLOCATING PROCESS B (417k)

- BEGIN FROM PARTITION 1 (388k): INSUFFICIENT.
- MOVE TO PARTITION 2 (700k): $700k \geq 417k$ ☒ ALLOCATE HERE.
- REMAINING SIZE: $700k - 417k = 283k$.
- UPDATED PARTITIONS:
- PARTITION 1: 388k

- PARTITION 2: 283k (REMAINING FREE)
- PARTITION 3: 300k
- PARTITION 4: 200k

STEP 3: ALLOCATING PROCESS C (112k)

- START AT PARTITION 1 (388k): SUFFICIENT ☒ ALLOCATE.
- REMAINING SIZE: $388k - 112k = 276k$.
- UPDATED PARTITIONS:
- PARTITION 1: 276k
- PARTITION 2: 283k
- PARTITION 3: 300k
- PARTITION 4: 200k

STEP 4: ALLOCATING PROCESS D (426k)

- START AT PARTITION 1 (276k): INSUFFICIENT.
- PARTITION 2 (283k): INSUFFICIENT.
- PARTITION 3 (300k): INSUFFICIENT.
- PARTITION 4 (200k): INSUFFICIENT.
- NO SUITABLE PARTITION AVAILABLE, SO PROCESS D CANNOT BE ALLOCATED WITH FIRST-FIT.

MEMORY ALLOCATION USING BEST-FIT ALGORITHM

NOW, LET'S SEE HOW BEST-FIT WOULD ALLOCATE THESE PROCESSES.

STEP 1: ALLOCATING PROCESS A (212k)

- SEARCH ALL FREE PARTITIONS:
- 600k, 700k, 300k, 200k
- THE SMALLEST PARTITION THAT FITS 212k:
- PARTITION 3: 300k (FITS, AND IS THE SMALLEST SUITABLE)
- ALLOCATE HERE.
- REMAINING SIZE: $300k - 212k = 88k$.
- UPDATED PARTITIONS:
- PARTITION 1: 600k
- PARTITION 2: 700k
- PARTITION 3: 88k (REMAINING FREE)
- PARTITION 4: 200k

STEP 2: ALLOCATING PROCESS B (417k)

- SEARCH ALL FREE PARTITIONS:
- 600k, 700k, 88k, 200k
- SUITABLE PARTITIONS:
- 600k AND 700k
- SMALLEST SUITABLE: 600k
- ALLOCATE HERE.
- REMAINING SIZE: $600k - 417k = 183k$.
- UPDATED PARTITIONS:

- PARTITION 1: 183k
- PARTITION 2: 700k
- PARTITION 3: 88k
- PARTITION 4: 200k

STEP 3: ALLOCATING PROCESS C (112k)

- SEARCH ALL FREE PARTITIONS:
- 183k, 700k, 88k, 200k
- SUITABLE PARTITIONS:
- 183k AND 200k
- SMALLEST SUITABLE: 183k
- ALLOCATE HERE.
- REMAINING SIZE: $183k - 112k = 71k$.
- UPDATED PARTITIONS:
- PARTITION 1: 71k
- PARTITION 2: 700k
- PARTITION 3: 88k
- PARTITION 4: 200k

STEP 4: ALLOCATING PROCESS D (426k)

- SEARCH ALL FREE PARTITIONS:
- 71k, 700k, 88k, 200k
- NONE ARE LARGE ENOUGH TO FIT 426k.
- PROCESS D CANNOT BE ALLOCATED WITH BEST-FIT.

COMPARISON AND ANALYSIS OF ALLOCATION OUTCOMES

FROM THE ABOVE STEPS, WE OBSERVE:

PROCESS	FIRST-FIT ALLOCATION	BEST-FIT ALLOCATION
A (212k)	PARTITION 1 (600k)	PARTITION 3 (300k)
B (417k)	PARTITION 2 (700k)	PARTITION 1 (600k)
C (112k)	PARTITION 1 (388k)	PARTITION 1 (183k)
D (426k)	NOT ALLOCATED	NOT ALLOCATED

KEY OBSERVATIONS:

- FIRST-FIT TENDS TO ALLOCATE PROCESSES QUICKLY BUT CAN LEAVE SMALL FRAGMENTS AT THE BEGINNING OF MEMORY, WHICH MAY PREVENT FUTURE ALLOCATIONS.
- BEST-FIT AIMS TO MINIMIZE WASTED SPACE BY CHOOSING THE SMALLEST SUITABLE PARTITION BUT MAY CAUSE EXTERNAL FRAGMENTATION OVER TIME.
- BOTH ALGORITHMS FAILED TO ALLOCATE PROCESS D (426k) IN THIS SCENARIO, ILLUSTRATING THE LIMITATION WHEN LARGE CONTIGUOUS FREE MEMORY IS UNAVAILABLE.

ADVANTAGES AND DISADVANTAGES OF EACH ALGORITHM IN PRACTICE

FIRST-FIT

ADVANTAGES:

- FASTER ALLOCATION DUE TO EARLY STOPPING ONCE A SUITABLE PARTITION IS FOUND.
- SIMPLE TO IMPLEMENT AND EFFICIENT FOR SYSTEMS WITH FREQUENT PROCESS ARRIVALS.

DISADVANTAGES:

- CAN CAUSE INTERNAL FRAGMENTATION AT THE BEGINNING OF MEMORY.
- MAY LEAD TO INEFFICIENT USE OF MEMORY OVER TIME, ESPECIALLY WITH MANY SMALL LEFTOVER PARTITIONS.

BEST-FIT

ADVANTAGES:

- MORE EFFICIENT UTILIZATION OF MEMORY BY MINIMIZING INTERNAL FRAGMENTATION.
- BETTER SUITED FOR SYSTEMS WHERE MEMORY WASTE REDUCTION IS CRITICAL.

DISADVANTAGES:

- SLOWER DUE TO THE NEED TO SEARCH THE ENTIRE LIST FOR THE BEST FIT.
- CAN LEAD TO EXTERNAL FRAGMENTATION, MAKING IT DIFFICULT TO FIND LARGE CONTIGUOUS FREE SPACES OVER TIME.

CONCLUSION

IN SUMMARY, BOTH FIRST-FIT AND BEST-FIT ALGORITHMS HAVE THEIR UNIQUE STRENGTHS AND WEAKNESSES. WHEN GIVEN THE SPECIFIC MEMORY PARTITIONS OF 600k, 700k, 300k, AND 200k IN ORDER, THEIR ALLOCATION BEHAVIORS DIFFER. FIRST-FIT ALLOCATES PROCESSES QUICKLY BY CHOOSING THE FIRST SUITABLE PARTITION, WHICH SOMETIMES LEADS TO INEFFICIENT MEMORY UTILIZATION AND FRAGMENTATION. BEST-FIT, ON THE OTHER HAND,

FREQUENTLY ASKED QUESTIONS

HOW DOES THE FIRST-FIT MEMORY ALLOCATION ALGORITHM ALLOCATE MEMORY FOR PARTITIONS OF 600k, 700k, 300k, AND 200k IN ORDER?

FIRST-FIT ALLOCATES EACH PROCESS TO THE FIRST AVAILABLE PARTITION THAT IS LARGE ENOUGH. STARTING WITH THE 600k PARTITION, IT ASSIGNS PROCESSES SEQUENTIALLY, FILLING THE EARLIEST SUITABLE PARTITION BEFORE MOVING TO THE NEXT.

HOW DOES THE BEST-FIT MEMORY ALLOCATION ALGORITHM ALLOCATE MEMORY FOR THE PARTITIONS 600k, 700k, 300k, AND 200k IN ORDER?

BEST-FIT SEARCHES THE ENTIRE LIST OF PARTITIONS TO FIND THE SMALLEST PARTITION THAT IS LARGE ENOUGH FOR EACH PROCESS, ALLOCATING PROCESSES TO THE MOST SUITABLE (SMALLEST SUFFICIENT) PARTITION IN ORDER.

WHAT IS THE MAIN DIFFERENCE BETWEEN FIRST-FIT AND BEST-FIT WHEN ALLOCATING MEMORY WITH PARTITIONS 600k, 700k, 300k, 200k?

FIRST-FIT ALLOCATES PROCESSES TO THE FIRST SUITABLE PARTITION ENCOUNTERED, WHILE BEST-FIT SEARCHES ALL PARTITIONS TO FIND THE SMALLEST SUITABLE ONE, AIMING TO REDUCE WASTED SPACE.

GIVEN THE MEMORY PARTITIONS 600k, 700k, 300k, 200k, HOW DOES FIRST-FIT ALLOCATE A PROCESS REQUIRING 250k?

FIRST-FIT ALLOCATES THE PROCESS TO THE FIRST PARTITION LARGE ENOUGH, WHICH IS 600k IN THIS CASE, SO IT ASSIGNS THE PROCESS THERE, LEAVING REMAINING SPACE OF 350k.

GIVEN THE SAME PARTITIONS, HOW WOULD BEST-FIT ALLOCATE A PROCESS REQUIRING 250k?

BEST-FIT SEARCHES ALL PARTITIONS AND ALLOCATES THE PROCESS TO THE SMALLEST PARTITION THAT CAN FIT 250k, WHICH IS 300k, LEAVING 50k UNUSED.

IF A PROCESS REQUIRES 150k, HOW DO FIRST-FIT AND BEST-FIT ALLOCATE MEMORY IN THE GIVEN PARTITIONS?

BOTH ALGORITHMS WILL ALLOCATE TO THE FIRST SUITABLE PARTITION. FIRST-FIT ASSIGNS TO 600k (REMAINING 450k), WHILE BEST-FIT FINDS 300k (REMAINING 150k) AS THE BEST FIT.

HOW DOES MEMORY UTILIZATION DIFFER BETWEEN FIRST-FIT AND BEST-FIT FOR THE GIVEN PARTITIONS?

FIRST-FIT TENDS TO PRODUCE LARGER LEFTOVER FRAGMENTS BY ALLOCATING TO THE EARLIEST SUITABLE PARTITIONS, WHILE BEST-FIT AIMS TO MINIMIZE WASTED SPACE BY CHOOSING THE SMALLEST ADEQUATE PARTITION, OFTEN LEADING TO BETTER UTILIZATION.

WHAT HAPPENS IF A PROCESS REQUIRING 650k IS ALLOCATED USING FIRST-FIT AND BEST-FIT?

USING FIRST-FIT, IT WOULD ALLOCATE TO THE 700k PARTITION, LEAVING 50k UNUSED. BEST-FIT WOULD ALSO ALLOCATE TO 700k, AS IT'S THE SMALLEST PARTITION LARGE ENOUGH, RESULTING IN SIMILAR LEFTOVER SPACE.

IN THE CONTEXT OF THE GIVEN PARTITIONS, WHICH ALGORITHM TENDS TO LEAVE SMALLER RESIDUAL FRAGMENTS, FIRST-FIT OR BEST-FIT?

BEST-FIT GENERALLY LEAVES SMALLER RESIDUAL FRAGMENTS BECAUSE IT ALLOCATES PROCESSES TO THE SMALLEST SUITABLE PARTITIONS, REDUCING WASTAGE.

HOW WOULD THE SEQUENCE OF ALLOCATIONS DIFFER BETWEEN FIRST-FIT AND BEST-FIT FOR MULTIPLE PROCESSES ARRIVING SEQUENTIALLY?

FIRST-FIT ALLOCATES EACH PROCESS TO THE FIRST AVAILABLE SUITABLE PARTITION, POTENTIALLY CAUSING FRAGMENTATION. BEST-FIT SEARCHES FOR THE BEST MATCH EACH TIME, OFTEN LEADING TO MORE EFFICIENT SPACE UTILIZATION AND LESS FRAGMENTATION OVER TIME.

ADDITIONAL RESOURCES

MEMORY ALLOCATION STRATEGIES: AN EXPERT ANALYSIS OF FIRST-FIT AND BEST-FIT WITH GIVEN PARTITIONS

MEMORY MANAGEMENT IS A FOUNDATIONAL ASPECT OF OPERATING SYSTEMS, DIRECTLY IMPACTING SYSTEM PERFORMANCE, EFFICIENCY, AND RESPONSIVENESS. AMONG THE VARIOUS STRATEGIES EMPLOYED, FIRST-FIT AND BEST-FIT ARE TWO OF THE MOST WIDELY USED ALGORITHMS FOR DYNAMIC MEMORY ALLOCATION. TO UNDERSTAND THEIR PRACTICAL IMPLICATIONS, LET'S

EXAMINE A DETAILED SCENARIO INVOLVING SPECIFIC MEMORY PARTITIONS AND ANALYZE HOW EACH STRATEGY MANAGES INCOMING PROCESSES.

UNDERSTANDING THE SCENARIO: THE GIVEN MEMORY PARTITIONS

SUPPOSE WE HAVE A SEQUENCE OF MEMORY PARTITIONS ARRANGED IN ORDER, EACH WITH A SPECIFIC SIZE:

- PARTITION 1: 600 KB
- PARTITION 2: 700 KB
- PARTITION 3: 300 KB
- PARTITION 4: 200 KB

THIS ORDERED SET OF PARTITIONS IS CRUCIAL BECAUSE BOTH FIRST-FIT AND BEST-FIT ALGORITHMS MAKE ALLOCATION DECISIONS BASED ON THE CURRENT STATE OF MEMORY PARTITIONS, WHICH CAN CHANGE OVER TIME AS PROCESSES ARE ALLOCATED AND DEALLOCATED.

MEMORY ALLOCATION STRATEGIES: AN OVERVIEW

BEFORE DIVING INTO HOW EACH ALGORITHM PERFORMS WITH THE GIVEN PARTITIONS, IT'S ESSENTIAL TO UNDERSTAND THEIR FUNDAMENTAL PRINCIPLES.

FIRST-FIT ALGORITHM

DEFINITION: THE FIRST-FIT ALGORITHM ALLOCATES THE FIRST AVAILABLE PARTITION THAT IS LARGE ENOUGH TO ACCOMMODATE THE PROCESS. IT SCANS THE LIST OF MEMORY BLOCKS SEQUENTIALLY FROM THE BEGINNING AND CHOOSES THE FIRST SUITABLE ONE.

ADVANTAGES:

- FASTER IN EXECUTION BECAUSE IT DOESN'T SEARCH THE ENTIRE LIST IF A SUITABLE PARTITION IS FOUND EARLY.
- SIMPLE TO IMPLEMENT.

DISADVANTAGES:

- CAN LEAD TO FRAGMENTATION, AS SMALL LEFTOVER SPACES MIGHT BE UNUSABLE LATER.
- TENDS TO ALLOCATE TO THE EARLIEST SUITABLE PARTITION, WHICH MIGHT CAUSE INEFFICIENT UTILIZATION OVER TIME.

BEST-FIT ALGORITHM

DEFINITION: THE BEST-FIT ALGORITHM SEARCHES THE ENTIRE LIST OF PARTITIONS AND ALLOCATES THE PROCESS TO THE SMALLEST AVAILABLE PARTITION THAT IS BIG ENOUGH TO HOLD IT. THIS APPROACH AIMS TO MINIMIZE WASTED SPACE.

ADVANTAGES:

- TENDS TO LEAVE SMALLER LEFTOVER PARTITIONS, REDUCING INTERNAL FRAGMENTATION.
- CAN POTENTIALLY MAKE BETTER USE OF MEMORY.

DISADVANTAGES:

- SLIGHTLY SLOWER DUE TO THE NEED TO SEARCH THE ENTIRE LIST FOR THE BEST FIT.

- MAY RESULT IN SMALL FRAGMENTS THAT ARE TOO SMALL FOR FUTURE ALLOCATIONS.

APPLYING FIRST-FIT AND BEST-FIT TO THE GIVEN PARTITIONS

LET'S ASSUME WE HAVE A SET OF PROCESSES ARRIVING SEQUENTIALLY, EACH REQUIRING A CERTAIN AMOUNT OF MEMORY. TO ANALYZE THE ALGORITHMS, WE NEED TO DEFINE THESE PROCESSES.

PROCESS SEQUENCE (EXAMPLE):

1. P1: 100 KB
2. P2: 200 KB
3. P3: 300 KB
4. P4: 600 KB
5. P5: 250 KB

NOTE: THE SEQUENCE AND SIZES ARE CHOSEN TO ILLUSTRATE HOW THE ALGORITHMS BEHAVE WITH VARYING PROCESS SIZES. ACTUAL BEHAVIOR CAN VARY BASED ON PROCESS ORDER.

FIRST-FIT ALLOCATION ANALYSIS

STEP-BY-STEP PROCESS:

PROCESS P1 (100 KB):

- START FROM PARTITION 1 (600 KB): $600\text{ KB} \geq 100\text{ KB}$ ☑ ALLOCATE HERE.
- REMAINING PARTITIONS:
- PARTITION 1: $600\text{ KB} - 100\text{ KB} = 500\text{ KB}$ (ALLOCATED)
- PARTITIONS 2, 3, 4 REMAIN UNCHANGED.

PROCESS P2 (200 KB):

- START FROM PARTITION 1: $500\text{ KB} \geq 200\text{ KB}$ ☑ ALLOCATE HERE.
- REMAINING:
- PARTITION 1: $500\text{ KB} - 200\text{ KB} = 300\text{ KB}$

PROCESS P3 (300 KB):

- START FROM PARTITION 1: $300\text{ KB} \geq 300\text{ KB}$ ☑ ALLOCATE HERE.
- REMAINING:
- PARTITION 1: $300\text{ KB} - 300\text{ KB} = 0\text{ KB}$ (FULLY ALLOCATED)
- NEXT PARTITIONS ARE NOT CONSIDERED AS THE PROCESS IS ALLOCATED.

PROCESS P4 (600 KB):

- PARTITION 1: 0 KB ☑ NOT SUITABLE.
- PARTITION 2 (700 KB): $700\text{ KB} \geq 600\text{ KB}$ ☑ ALLOCATE HERE.
- REMAINING:
- PARTITION 2: $700\text{ KB} - 600\text{ KB} = 100\text{ KB}$

PROCESS P5 (250 KB):

- PARTITION 2: 100 KB ☑ TOO SMALL.
- PARTITION 3: $300\text{ KB} \geq 250\text{ KB}$ ☑ ALLOCATE HERE.
- REMAINING:
- PARTITION 3: $300\text{ KB} - 250\text{ KB} = 50\text{ KB}$

SUMMARY OF FIRST-FIT ALLOCATIONS:

PROCESS	ALLOCATED PARTITION	PARTITION REMAINING SPACE
P1	PARTITION 1 (600 KB)	500 KB (AFTER ALLOCATION)
P2	PARTITION 1	300 KB
P3	PARTITION 1	0 KB
P4	PARTITION 2	100 KB
P5	PARTITION 3	50 KB

BEST-FIT ALLOCATION ANALYSIS

STEP-BY-STEP PROCESS:

PROCESS P1 (100 KB):

- SEARCH ALL PARTITIONS:
- PARTITION 1: 600 KB
- PARTITION 2: 700 KB
- PARTITION 3: 300 KB
- PARTITION 4: 200 KB
- SMALLEST SUITABLE PARTITION: PARTITION 4 (200 KB)
- ALLOCATE HERE:
- PARTITION 4: $200 \text{ KB} - 100 \text{ KB} = 100 \text{ KB}$ REMAINING

PROCESS P2 (200 KB):

- SEARCH REMAINING PARTITIONS:
- PARTITION 1: 600 KB
- PARTITION 2: 700 KB
- PARTITION 3: 300 KB
- PARTITION 4: 100 KB (NOT SUITABLE)
- SMALLEST SUITABLE: PARTITION 3 (300 KB)
- ALLOCATE:
- PARTITION 3: $300 \text{ KB} - 200 \text{ KB} = 100 \text{ KB}$ REMAINING

PROCESS P3 (300 KB):

- SEARCH:
- PARTITION 1: 600 KB
- PARTITION 2: 700 KB
- PARTITION 3: 100 KB (NOT SUITABLE)
- PARTITION 4: 100 KB (NOT SUITABLE)
- SMALLEST SUITABLE: PARTITION 1 (600 KB)
- ALLOCATE:
- PARTITION 1: $600 \text{ KB} - 300 \text{ KB} = 300 \text{ KB}$ REMAINING

PROCESS P4 (600 KB):

- SEARCH:
- PARTITION 1: 300 KB (NOT SUITABLE)
- PARTITION 2: 700 KB [?] SUITABLE
- PARTITION 3: 100 KB
- PARTITION 4: 100 KB
- SMALLEST SUITABLE: PARTITION 2 (700 KB)
- ALLOCATE:
- PARTITION 2: $700 \text{ KB} - 600 \text{ KB} = 100 \text{ KB}$ REMAINING

PROCESS P5 (250 KB):

- SEARCH:

- PARTITION 1: 300 KB ☐ SUITABLE
- PARTITION 2: 100 KB (NOT SUITABLE)
- PARTITION 3: 100 KB
- PARTITION 4: 100 KB
- SMALLEST SUITABLE: PARTITION 1 (300 KB)
- ALLOCATE:
- PARTITION 1: 300 KB - 250 KB = 50 KB REMAINING

SUMMARY OF BEST-FIT ALLOCATIONS:

PROCESS	ALLOCATED PARTITION	PARTITION REMAINING SPACE
P1	PARTITION 4 (200 KB)	100 KB
P2	PARTITION 3 (300 KB)	100 KB
P3	PARTITION 1 (600 KB)	300 KB
P4	PARTITION 2 (700 KB)	100 KB
P5	PARTITION 1	50 KB

COMPARISON AND INSIGHTS

EFFICIENCY AND FRAGMENTATION:

- FIRST-FIT TENDS TO ALLOCATE PROCESSES IN THE EARLIEST SUITABLE PARTITION, WHICH CAN CAUSE THE “FILLING FROM THE FRONT” ISSUE. THIS OFTEN RESULTS IN HIGHER EXTERNAL FRAGMENTATION OVER TIME BECAUSE SMALL LEFTOVER SPACES AT THE BEGINNING OF THE MEMORY LIST REMAIN UNUSED, EVEN IF LARGER PROCESSES COULD FIT INTO OTHER SCATTERED FREE BLOCKS.
- BEST-FIT AIMS TO REDUCE WASTAGE BY ALLOCATING TO THE SMALLEST SUITABLE PARTITION. WHILE THIS CAN MINIMIZE INTERNAL FRAGMENTATION AND MAKE EFFICIENT USE OF MEMORY BLOCKS, IT OFTEN LEADS TO INCREASED SEARCH TIMES DUE TO SCANNING THE ENTIRE LIST FOR THE BEST FIT. ALSO, IT CAN CREATE NUMEROUS SMALL FRAGMENTS THAT ARE TOO SMALL FOR FUTURE ALLOCATIONS, LEADING TO EXTERNAL FRAGMENTATION OVER TIME.

IN OUR SCENARIO:

- FIRST-FIT ALLOCATED PROCESSES TO THE EARLIEST SUITABLE BLOCKS, LEADING TO A RELATIVELY STRAIGHTFORWARD ALLOCATION PATTERN BUT POTENTIALLY LEAVING LARGER SPACES UNUSED AT THE END.
- BEST-FIT MADE MORE NUANCED CHOICES, OFTEN ALLOCATING TO SMALLER PARTITIONS TO PRESERVE LARGER ONES FOR BIGGER PROCESSES. THIS RESULTED IN A MORE BALANCED UTILIZATION OF THE MEMORY PARTITIONS BUT COULD ALSO LEAD TO FRAGMENTATION IF PROCESSES ARE DEALLOCATED AND REALLOCATED REPEATEDLY.

IMPLICATIONS IN REAL-WORLD SYSTEMS

UNDERSTANDING HOW THESE ALGORITHMS BEHAVE WITH SPECIFIC PARTITION SIZES IS CRUCIAL FOR SYSTEM DESIGNERS AND DEVELOPERS:

- FIRST-FIT IS FAVORED FOR ITS SPEED AND SIMPLICITY, ESPECIALLY IN SYSTEMS WHERE ALLOCATION SPEED IS CRITICAL, AND FRAGMENTATION CAN BE TOLERATED OR MANAGED.
- BEST-FIT IS CHOSEN WHEN MEMORY UTILIZATION EFFICIENCY IS PRIORITIZED, AND THE SYSTEM CAN HANDLE THE OVERHEAD OF SEARCHING FOR THE OPTIMAL PARTITION.

PRACTICAL RECOMMENDATIONS:

- FOR SYSTEMS WITH PREDICTABLE WORKLOADS AND WHERE FRAGMENTATION IS MANAGEABLE, BEST-FIT CAN MAXIMIZE MEMORY USE.
- FOR REAL-TIME SYSTEMS OR THOSE REQUIRING RAPID ALLOCATION, FIRST-FIT'S SPEED ADVANTAGES OUTWEIGH THE POTENTIAL

Given The Memory Partitions Of 600k 700k 300k 200k In Order How Would Each Of The First Fit And Best Fit

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